

Some Remarks on the Generalized Scheme of Reduction to Absurdity and Generalized Hypothetical Syllogism in Fuzzy Logic

Katarzyna Miś and Michał Baczyński and Piotr Helbin

Institute of Mathematics, University of Silesia in Katowice,

40-007 Katowice, Bankowa 14, Poland,

{kmis,michal.baczynski,piotr.helbin}@us.edu.pl

Abstract

In this paper we investigate two generalizations, in fuzzy logic, of classical scheme of reduction to absurdity. We compare them with two possible generalizations of classical hypothetical syllogism (in fuzzy logic) and we show that generalized hypothetical syllogism is more general. We present new results concerning solutions of an inequality and an equation connected directly with generalization of scheme of reduction to absurdity in fuzzy logic.

Keywords: Generalized reduction to absurdity, generalized hypothetical syllogism, fuzzy implications, R-implications, t-norms, fuzzy negations

1 Introduction

There are many reasoning schemas (rules of inferences) in classical logic, like modus (ponendo) ponens, modus (tollendo) tollens, scheme of disjunctive reasoning, law of contraposition, etc. They are also applied in the terms of fuzzy logic. Namely, they are used in approximate reasoning and/or fuzzy control. Recently, we have investigated generalized hypothetical syllogism in fuzzy logic [2] (see also [7]). This notion can be introduced from a T -transitivity in the following way

$$T(I(x, z), I(z, y)) \leq I(x, y), \quad x, y, z \in [0, 1], \quad (\text{HS})$$

where T is a t-norm and I a fuzzy implication. However, involving Zadeh's compositional rule of inference (CRI) [8] we can receive the following functional equation, satisfied for all $x, y \in [0, 1]$,

$$\sup_{z \in [0, 1]} (T(I(x, z), I(z, y))) = I(x, y). \quad (\text{GHS})$$

In this paper we investigate different scheme – reduction to absurdity (in Latin “*reductio ad absurdum*”). In general we can write it in fuzzy logic as follows

$$\begin{array}{ll} \text{RULE:} & \text{IF } x \text{ is not } A, \text{ THEN } y \text{ is } B \\ \text{FACT:} & y \text{ is not } B \\ \hline \text{CONCLUSION:} & x \text{ is } A \end{array}$$

where A, B are fuzzy sets that represent some properties. Based on rules from the Boolean algebra and important investigations from [6], where some generalizations of classical schemes of reasoning were examined in fuzzy logic, we can write the following inequality, which corresponds with the reduction to absurdity,

$$T(I(N(x), y), N(y)) \leq x, \quad x, y \in [0, 1]. \quad (\text{RA})$$

However, again using CRI it is possible to obtain the functional equation of the form

$$\sup_{y \in [0, 1]} T(I(N(x), y), N(y)) = x, \quad x \in [0, 1], \quad (\text{GRA})$$

where N is a fuzzy negation, T a t-norm and I a fuzzy implication.

The main goal of this article is to compare written above inequalities and equations for generalized hypothetical syllogism and generalized scheme of reduction to absurdity in fuzzy logic. Moreover, we give some new results concerning particular solutions of (RA) and (GRA).

The paper is organised as follows. Section 2 contains some important facts and definitions used in the sequel, while in Section 3 we present some properties regarding (RA) and some solutions of (RA) and (GRA) for several families of fuzzy implications. We also present some new results concerning (GHS).

2 Preliminaries

To make this work more self-contained, we place some of basic definitions concerning fuzzy connectives here. Note that the family of all increasing bijections $\varphi: [0, 1] \rightarrow [0, 1]$ will be denoted by the symbol Φ .

Definition 2.1 (see [3, 5]). A function $T: [0, 1]^2 \rightarrow [0, 1]$ is called a **triangular norm (t-norm in short)**,

if it satisfies, for all $x, y, z \in [0, 1]$, the following conditions:

- (T1) $T(x, y) = T(y, x)$,
- (T2) $T(x, T(y, z)) = T(T(x, y), z)$,
- (T3) $T(x, y) \leq T(x, z)$ for $y \leq z$, i.e., $T(x, \cdot)$ is non-decreasing,
- (T4) $T(x, 1) = x$.

Definition 2.2 (see [5]). A function $S: [0, 1]^2 \rightarrow [0, 1]$ is called a **triangular conorm** (**t-conorm** in short), if it satisfies, for all $x, y, z \in [0, 1]$, the following conditions:

- (S1) $S(x, y) = S(y, x)$,
- (S2) $S(x, S(y, z)) = S(S(x, y), z)$,
- (S3) $S(x, y) \leq S(x, z)$ for $y \leq z$, i.e., $S(x, \cdot)$ is non-decreasing,
- (S4) $S(x, 0) = x$.

Example 2.3. 1. The Łukasiewicz t-norm,

$$T_{LK}(x, y) = \max(x + y - 1, 0), \quad x, y \in [0, 1].$$

2. The Łukasiewicz t-conorm,

$$S_{LK}(x, y) = \min(x + y, 1), \quad x, y \in [0, 1].$$

Definition 2.4 (see [1, 5]). A non-increasing function $N: [0, 1] \rightarrow [0, 1]$ is called a **fuzzy negation**, if $N(0) = 1$, $N(1) = 0$. Moreover, a fuzzy negation N is called

- (i) **strict**, if it is strictly decreasing and continuous,
- (ii) **strong**, if it is an involution, i.e., $N(N(x)) = x$ for all $x \in [0, 1]$.

Example 2.5. 1. The classical negation N_C is given by $N_C(x) = 1 - x$, for $x \in [0, 1]$.

2. The least negation N_{D1} is given by

$$N_{D1}(x) = \begin{cases} 1, & x = 0, \\ 0, & x > 0, \end{cases} \quad x \in [0, 1].$$

Definition 2.6 ([1, Definition 2.3.14]). Let T be a t-norm and N a fuzzy negation. We say that a pair (T, N) satisfies the law of contradiction if

$$T(x, N(x)) = 0, \quad x \in [0, 1]. \quad (\text{LC})$$

Now, we recall the definition and some important properties of fuzzy implications.

Definition 2.7 (see [1, 3]). A function $I: [0, 1]^2 \rightarrow [0, 1]$ is called a **fuzzy implication**, if it satisfies the following conditions.

- (I1) I is non-increasing with respect to the first variable,
- (I2) I is non-decreasing with respect to the second variable,
- (I3) $I(0, 0) = I(1, 1) = 1$ and $I(1, 0) = 0$.

The family of all fuzzy implications will be denoted by \mathcal{FI} .

Definition 2.8 (see [1]). We say that a fuzzy implication I satisfies

- (i) the **identity principle**, if

$$I(x, x) = 1, \quad x \in [0, 1], \quad (\text{IP})$$

- (ii) the **left neutrality property**, if

$$I(1, y) = y, \quad y \in [0, 1], \quad (\text{NP})$$

- (iii) the **ordering property**, if

$$x \leq y \iff I(x, y) = 1, \quad x, y \in [0, 1]. \quad (\text{OP})$$

- (iv) **law of contraposition with respect to N** , if

$$I(x, y) = I(N(y), N(x)), \quad x, y \in [0, 1], \quad (\text{CP})$$

- (v) **law of left contraposition with respect to N** , if

$$I(N(x), y) = I(N(y), x), \quad x, y \in [0, 1], \quad (\text{L-CP})$$

- (vi) **law of right contraposition with respect to N** , if

$$I(x, N(y)) = I(y, N(x)), \quad x, y \in [0, 1]. \quad (\text{R-CP})$$

Definition 2.9 ([1, Definition 1.4.15]). Let I be a fuzzy implication. A function $N_I: [0, 1] \rightarrow [0, 1]$ given by

$$N_I(x) = I(x, 0), \quad x \in [0, 1],$$

is called the natural negation of I .

Let us also recall definitions of two families of fuzzy implications.

Definition 2.10 ([1, Definition 2.5.1]). A function $I: [0, 1]^2 \rightarrow [0, 1]$ is called a residual implication (R-implication for short) if there exists a t-norm T such that

$$I(x, y) = \sup\{t \in [0, 1] \mid T(x, t) \leq y\}, \quad x, y \in [0, 1].$$

If I is generated from a t-norm T , then it will be denoted by I_T .

Definition 2.11 ([1, Definition 2.4.1]). A function $I: [0, 1]^2 \rightarrow [0, 1]$ is called an (S, N) -implication, if there exist a t-conorm S and a fuzzy negation N such that

$$I(x, y) = S(N(x), y), \quad x, y \in [0, 1].$$

If I is generated from a t-conorm S and a fuzzy negation N , then it will be denoted by $I_{S,N}$.

3 Properties of (RA) and (GRA)

Let us start with some general properties of triplets (T, I, N) satisfying (RA).

Proposition 3.1. *Let $I \in \mathcal{FI}$, T be a t-norm and N be a fuzzy negation. Next, let the triplet (T, I, N) satisfies (RA).*

1. *If I satisfies (NP), then (T, N) satisfies (LC).*
2. *If N_I is injective, then $N_I^{-1} \leq N$.*
3. *If I satisfies (NP), then $N \leq N_T$.*
4. *If $N \neq N_{D_1}$ and I satisfies (NP), then T has zero-divisors ($a \in (0, 1)$ is a zero-divisor if there exists $b \in (0, 1)$ such that $T(a, b) = 0$).*
5. *If $N_I(x) \neq 0$, $x \in [0, 1]$, T is continuous and I satisfies (NP), then there exists $\varphi \in \Phi$ such that*

$$\begin{aligned} T(x, y) &= (T_{\mathbf{LK}})_\varphi(x, y) \\ &= \varphi^{-1}(\max(0, \varphi(x) + \varphi(y) - 1)), \end{aligned}$$

for all $x, y \in [0, 1]$.

6. *If $I_1 \in \mathcal{FI}$, T_2 is a t-norm such that $T_2 \leq T$ and $I_1 \leq I$, then (T_2, I_1) satisfies (RA).*

Proof. 1. It is enough to take $x = 0$. Hence we have $0 \leq T(N(y), I(1, y)) = T(N(y), y) \leq 0$, for all $y \in [0, 1]$.

2. If we take $y = 0$ and we assume N_I is injective, then for every $x \in [0, 1]$ we obtain

$$\begin{aligned} T(1, I(N(x), 0)) &\leq x, \\ I(N(x), 0) &\leq x, \\ N_I(N(x)) &\leq x, \\ N(x) &\geq N_I^{-1}(x). \end{aligned}$$

3. It is immediate from the point 1. and from the formula of N_T (negation induced by T) given by

$$N_T(x) = \sup\{y \in [0, 1] \mid T(x, y) = 0\},$$

for $x \in [0, 1]$.

4. If $N \neq N_{D_1}$, then there exists $x \in (0, 1)$ such that $N(x) > 0$. Hence from the point 1. T has zero-divisors because (T, N) satisfies (LC).

5. If $N_I(x) > 0$ for $x \in [0, 1]$, then every such x is a zero-divisor of T . Moreover, the only continuous t-norm such the set of zero-divisors is $(0, 1)$ is $T = (T_{\mathbf{LK}})_\varphi$, for some $\varphi \in \Phi$ (see [5, Remark 2.4, Proposition 5.10]).

6. It is straightforward from the following inequalities,

$$\begin{aligned} T_2(N(y), I_1(N(x), y)) &\leq T_2(N(y), I(N(x), y)) \\ &\leq T(N(y), I(N(x), y)) \\ &\leq x, \end{aligned}$$

satisfied for all $x, y \in [0, 1]$. □

Now, we recall one result concerning triplets (T, I, N) satisfying the (RA) but with some strong assumptions regarding T and N .

Theorem 3.2 ([6, Proposition 6.3]). *Let T be a continuous t-norm, N a strong negation and let $I \in \mathcal{FI}$ satisfy (NP). Then the following statements are equivalent.*

- (i) *The triplet (T, I, N) satisfies (RA).*
- (ii) *there exists $\varphi \in \Phi$ such that $T = (T_{\mathbf{LK}})_\varphi$, $N \leq (N_{\mathbf{C}})_\varphi$ and*

$$\varphi(I(x, y)) \leq 1 - \varphi(N(y)) + \varphi(N(x)), \quad x, y \in [0, 1].$$

Of course we can find some solutions of (RA), where T is non-continuous.

Example 3.3. Let us consider the drastic (non-continuous) t-norm given by

$$T_{\mathbf{D}}(x, y) = \begin{cases} 0, & x, y \in [0, 1), \\ \min(x, y), & \text{otherwise,} \end{cases}$$

and the Gödel implication given by

$$I_{\mathbf{GD}}(x, y) = \begin{cases} 1, & x \leq y, \\ y, & x > y. \end{cases}$$

Then it can be quite easily verified that the triplet $(T_{\mathbf{D}}, I_{\mathbf{GD}}, N_{\mathbf{C}})$ satisfies (RA).

It is obvious that if a triplet (T, I, N) satisfies (GRA), then such triplet satisfies also (RA). Similarly, if a pair (T, I) satisfies (GHS), then (T, I) satisfies also (HS). In general, without additional assumptions, the opposite implications are not true.

Example 3.4. The triplet $(T_{\mathbf{D}}, I_{\mathbf{GD}}, N_{\mathbf{C}})$ satisfies (RA) and does not satisfy (GRA). Furthermore,

the pair (T_D, I_{RC}) satisfies (HS) and does not satisfy (GHS), where I_{RC} is the Reichenbach implication given by the formula $I_{RC}(x, y) = 1 - x + xy$, for $x, y \in [0, 1]$.

The next results show some sufficient conditions to satisfy (GHS) and (GRA).

Proposition 3.5. *Let T be a t-norm and $I \in \mathcal{FI}$. If the pair (T, I) satisfies (HS) and I satisfies (IP), then (T, I) satisfies (GHS).*

Proof. Let $x, y \in [0, 1]$. From one side, from (HS) and monotonicity of supremum we have

$$\sup_{z \in [0, 1]} (T(I(x, z), I(z, y))) \leq \sup_{z \in [0, 1]} I(x, y) = I(x, y).$$

On the other side,

$$\begin{aligned} \sup_{z \in [0, 1]} (T(I(x, z), I(z, y))) &\geq T(I(x, y), I(y, y)) \\ &= T(I(x, y), 1) = I(x, y). \end{aligned}$$

Therefore (GHS) is true for the pair (T, I) . \square

However, the above condition is not necessary.

Example 3.6. Let I_{KD} be the Kleene-Dienes implication given by

$$I_{KD}(x, y) = \max(1 - x, y), \quad x, y \in [0, 1].$$

Then the pair (T_D, I_{KD}) satisfies (HS) although I_{KD} does not satisfy (IP). Moreover, this pair satisfies (GHS). Indeed, if $x = 0$, then for all $y \in [0, 1]$

$$\begin{aligned} \sup_{z \in [0, 1]} T_D(I_{KD}(0, z), I_{KD}(z, y)) &= \sup_{z \in [0, 1]} T_D(1, \max(1 - z, y)) \\ &= \sup_{z \in [0, 1]} \max(1 - z, y) = 1 = I_{KD}(0, y), \end{aligned}$$

and similarly, if $y = 1$, then for all $z \in [0, 1]$ we have

$$\begin{aligned} \sup_{z \in [0, 1]} T_D(I_{KD}(x, z), I_{KD}(z, 1)) &= \sup_{z \in [0, 1]} T_D(\max(1 - x, z), 1) \\ &= \sup_{z \in [0, 1]} \max(1 - x, z) = 1 = I_{KD}(x, 1). \end{aligned}$$

Next, from the definition of the drastic t-norm T_D , for every $x \in (0, 1]$, $y \in [0, 1)$ and $z \in (0, 1)$ we have

$$T_D(I_{KD}(x, z), I_{KD}(z, y)) = 0,$$

thus when $x \in (0, 1]$ and $y \in [0, 1)$ we have

$$\begin{aligned} \sup_{z \in [0, 1]} T_D(I_{KD}(x, z), I_{KD}(z, y)) &= \max(T_D(I_{KD}(x, 0), I_{KD}(0, y)), \\ &\quad T_D(I_{KD}(x, 1), I_{KD}(1, y))) \\ &= \max(1 - x, y) = I_{KD}(x, y). \end{aligned}$$

Therefore the pair (T_D, I_{KD}) satisfies (GHS).

Since each R-implication satisfies (IP) (see e.g. [1, Theorem 2.5.4]), the following fact is true for all R-implications.

Corollary 3.7. *For a t-norm T the following statements are equivalent.*

- (i) *The pair (T, I_T) satisfies (HS).*
- (ii) *The pair (T, I_T) satisfies (GHS).*

Now, let us return to (RA) and (GRA).

Proposition 3.8. *Let T be a t-norm, N be a strict negation and $I \in \mathcal{FI}$. If the triplet (T, I, N^{-1}) satisfies (RA) and $N_I \circ N^{-1} = id$ (i.e., $N_I = N$), then the triplet $(T, I, N^{-1}) = (T, I, N_I^{-1})$ satisfies (GRA).*

Proof. Let $x \in [0, 1]$. On the one hand, from (RA) and monotonicity of supremum we have

$$\sup_{y \in [0, 1]} T(I(N^{-1}(x), y), N^{-1}(y)) \leq \sup_{y \in [0, 1]} x = x.$$

On the other hand, from our assumption we obtain

$$\begin{aligned} \sup_{y \in [0, 1]} T(I(N^{-1}(x), y), N^{-1}(y)) &\geq T(I(N^{-1}(x), 0), N^{-1}(0)) \\ &= T(I(N^{-1}(x), 0), 1) = N_I \circ N(x) = x. \end{aligned}$$

Therefore (GRA) is true for the triplet (T, I, N^{-1}) . \square

Example 3.9. It can be easily checked that the triplet (T, I, N_I) , where $T = T_D$ and

$$I(x, y) = \max(1 - x^2, y), \quad x, y \in [0, 1],$$

satisfies (RA). From the above result we know that this triplet satisfies also (GRA).

Proposition 3.10. *Let T be a t-norm, $I \in \mathcal{FI}$ and let N_I be a strong negation.*

1. *If the pair (T, I) satisfies (GHS), then the triplet (T, I, N_I) satisfies (GRA).*
2. *If the pair (T, I) satisfies (HS), then the triplet (T, I, N_I) satisfies (RA).*

Proof. 1. Let us take $y = 0$ and substitute $N_I(x)$ instead of x in (GHS). Then we have the following equations

$$\begin{aligned} \sup_{z \in [0, 1]} T(I(N_I(x), z), I(z, 0)) &= I(N_I(x), 0), \\ \sup_{z \in [0, 1]} T(I(N_I(x), z), N_I(z)) &= N_I(N_I(x)), \\ \sup_{z \in [0, 1]} T(I(N_I(x), z), N_I(z)) &= x, \end{aligned}$$

which means that the triplet (T, I, N_I) satisfies (GRA).

2. This proof is similar to that above. \square

Therefore, in some cases we can apply the following theorem valid for R-implications.

Theorem 3.11 ([2, Theorem 4.12]). *Let T^* be a t -norm and T be a left-continuous t -norm. Then the following statements are equivalent.*

(i) *The pair (T^*, I_T) satisfies (GHS).*

(ii) $T^* \leq T$.

Corollary 3.12. *Let T^* be a t -norm, T be a left-continuous t -norm and N_{I_T} be a strong negation. Then the following statements are equivalent.*

(i) *The triplet (T^*, I_T, N_{I_T}) satisfies (GRA).*

(ii) $T^* \leq T$.

Proof. (i) \implies (ii) Since T is a left-continuous t -norm, the following equivalence is true (see [1, Proposition 2.5.2]), for all $x, y \in [0, 1]$,

$$T(x, y) \leq T(x, y) \iff I_T(x, T(x, y)) \geq y.$$

Let us recall that if $N = N_{I_T}$ is a strong negation, then I_T satisfies (L-CP) with N (see [1, Proposition 2.5.28]). Hence, for arbitrary fixed $x, y \in [0, 1]$, we have

$$\begin{aligned} T(x, y) &= \sup_{z \in [0, 1]} T^*(I_T(N(T(x, y))), z, N(z)) \\ &= \sup_{z \in [0, 1]} T^*(I_T(N(z), T(x, y)), N(z)) \\ &\geq T^*(I_T(N(N^{-1}(x)), T(x, y)), N(N^{-1}(x))) \\ &= T^*(I_T(x, T(x, y)), x) \geq T^*(x, y). \end{aligned}$$

(ii) \implies (i) If $T^* \leq T$, then from Theorem 3.11 we know that the pair (T^*, I_T) satisfies (GHS). Thus in virtue of Proposition 3.10 we obtain that the triplet (T^*, I_T, N_{I_T}) satisfies (GRA). \square

Moreover, we have also the following fact.

Proposition 3.13. *For a t -norm T the following statements are equivalent.*

(i) *The pair (T, I_T) satisfies (HS).*

(ii) T is left-continuous.

Proof. (i) \implies (ii) Suppose that the pair (T, I_T) satisfies (HS) and T is not left-continuous. Then there exist

$x, y, z \in [0, 1]$ such that $I_T(x, y) \geq z$ and $T(x, z) > y$ (see [1, Proposition 2.5.2]). Hence

$$\begin{aligned} y &< T(x, z) \leq T(x, I_T(x, y)) \\ &= T(I_T(1, x), I_T(x, y)) \leq I_T(1, y) \\ &= y; \end{aligned}$$

a contradiction.

(ii) \implies (i) Note that (HS) is nothing else but the T -transitivity – the property satisfied for pairs (T, I_T) , where T is a left-continuous t -norm (see [4, Proposition 1.6]). \square

Now we can formulate the following corollary.

Corollary 3.14. *Let T be a t -norm. If I_T satisfies (L-CP) with a negation N , then the following statements are equivalent.*

(i) *The triplet (T, I_T, N) satisfies (RA).*

(ii) T is left-continuous.

Proof. (i) \implies (ii) This part of the proof is similar to the proof of Proposition 3.13. Suppose that I_T is not left-continuous. Then there exist $x, y, z \in [0, 1]$ such that $I_T(y, x) \geq z$ and $T(y, z) > x$. Note that if I_T satisfies (L-CP) with a negation N , then $N = N_{I_T}$ is strong ([1, Proposition 2.5.26]), and therefore $y = N(y_0)$ for some $y_0 \in [0, 1]$. Hence

$$\begin{aligned} x &< T(y, z) \leq T(y, I(y, x)) \\ &= T(N(y_0), I_T(N(y_0), x)) \\ &= T(N(y_0), I_T(N(x), y_0)) \\ &\leq x; \end{aligned}$$

a contradiction.

(ii) \implies (i) If T is left-continuous and I_T satisfies (L-CP) with N , then again from [1, Proposition 2.5.2] we have

$$\begin{aligned} I_T(N(y), x) &\geq I_T(N(x), y) \iff \\ &T(N(y), I_T(N(x), y)) \leq x, \end{aligned}$$

for all $x, y \in [0, 1]$. \square

Remark 3.15. We know that if I_T satisfies (L-CP) with a negation N , then $N = N_{I_T}$ is strong (see [1, Proposition 2.5.26]). However it is not equivalent with left-continuity of T . Indeed, the Fodor implication given by

$$I_{\text{FD}}(x, y) = \begin{cases} 1, & x \leq y, \\ \max(1 - x, y), & x > y, \end{cases}$$

which satisfies (L-CP) with N_C can be generated from the non left-continuous t-norm T_{nM^*} given by

$$T_{nM^*}(x, y) = \begin{cases} 0, & x + y < 1, \\ \min(x, y), & \text{otherwise.} \end{cases}$$

One of sufficient conditions in such cases for T to be left-continuous can be satisfying (RA) by the triplet (T, I_T, N_{I_T}) .

Remark 3.16. Assumptions of having left-continuous t-norm and strong negation N were crucial in previous results. Note that for such fuzzy negation (S, N) -implication $I_{S,N}$ satisfies (L-CP) with N (see [1, Proposition 2.4.3]). However, if we consider such (S, N) -implications, which are also R-implications generated from a left-continuous t-norm, we obtain only $I_{S,N} = (I_{LK})_\varphi$, for $\varphi \in \Phi$, where I_{LK} is the Łukasiewicz implication given by the formula $I_{LK}(x, y) = \min(1 - x + y, 1)$, for all $x, y \in [0, 1]$.

Example 3.17. Despite this, we can find (S, N) -implications (which are not R-implications at the same time) satisfying (RA), for example the triplets (T_{LK}, I_{RC}, N_C) and (T_D, I_{RC}, N_C) satisfy (RA).

Theorem 3.18. Let T be a continuous t-norm, S be a t-conorm, N be a strong negation and let I be an (S, N) -implication. If the triplet (T, I, N) satisfies (GRA), then $T = (T_{LK})_\varphi$ and $N \leq (N_C)_\varphi$, for some $\varphi \in \Phi$. Moreover, if $T = (T_{LK})_\varphi$, $N \leq (N_C)_\varphi$ and $S \leq (S_{LK})_\varphi$, for some $\varphi \in \Phi$, then the triplet (T, I, N) satisfies (GRA).

Proof. Assume that the triplet (T, I, N) satisfies (GRA). From Theorem 3.2 we know that $T = (T_{LK})_\varphi$ and $N \leq (N_C)_\varphi$, for some $\varphi \in \Phi$.

Assume now that $T = (T_{LK})_\varphi$, $N \leq (N_C)_\varphi$ and $S \leq (S_{LK})_\varphi$, for some $\varphi \in \Phi$. Thus, for $x, y \in [0, 1]$, we obtain

$$\begin{aligned} T(N(y), I(N(x), y)) &= (T_{LK})_\varphi(N(y), S(x, y)) \\ &= \varphi^{-1}(\max(\varphi(N(y)) + \varphi(S(x, y)) - 1, 0)) \\ &\leq \varphi^{-1}(\max(\varphi((N_C)_\varphi(y)) + \varphi((S_{LK})_\varphi(x, y)) - 1, 0)) \\ &\leq \varphi^{-1}(\max(-\varphi(y) + \min(\varphi(x) + \varphi(y), 1), 0)) \\ &\leq \varphi^{-1}(\max(\min(\varphi(x), 1 - \varphi(y)), 0)) \\ &\leq \varphi^{-1}(\varphi(x)) = x, \end{aligned}$$

which proves that the triplet (T, I, N) satisfies (RA). However, $N_I = N$ and N is the strong negation. From Proposition 3.8 we obtain the thesis. \square

Let us finish with the following example which illustrates the last theorem.

Example 3.19. Consider the following functions

$$\begin{aligned} I(x, y) &= \max(1 - x, y), \quad x, y \in [0, 1], \\ T(x, y) &= \sqrt{\max(x^2 + y^2 - 1, 0)}, \quad x, y \in [0, 1], \\ N_1(x) &= \sqrt{1 - x^2}, \quad x \in [0, 1]. \end{aligned}$$

Thus, I is the Kleene-Dienes implication (it is an (S, N) -implication, cf. [1, Table 2.4]), $T = (T_{LK})_\varphi$ and $N_1 = (N_C)_\varphi$, where $\varphi(x) = x^2$, for all $x \in [0, 1]$. From the above theorem the triplet (T, I, N) satisfies (GRA), where $N = N_C \leq (N_C)_\varphi$.

4 Conclusions

We have investigated the scheme of reduction to absurdity (GRA) and generalized hypothetical syllogism (GHS). We presented some similar results for both of them. Also we shown that in some cases (GHS) is more general. Moreover, we presented conditions when (RA) is equivalent to (GRA).

Acknowledgement

M. Baczyński and K. Miś acknowledge the support of the National Science Centre, Poland, under Grant No. 2015/19/B/ST6/03259.

References

- [1] M. Baczyński, B. Jayaram, Fuzzy Implications, Vol. 231 of Studies in Fuzziness and Soft Computing, Springer, Berlin Heidelberg, 2008.
- [2] M. Baczyński, K. Miś, Selected properties of generalized hypothetical syllogism including the case of R-implications, in: M. J. et al. (eds) (Ed.), Information Processing and Management of Uncertainty in Knowledge-Based Systems. Theory and Foundations, Vol. 853 of Communications in Computer and Information Science, Springer, Cham, 2018, pp. 673–684.
- [3] J. Fodor, M. Roubens, Fuzzy Preference Modelling and Multicriteria Decision Support, Kluwer Academic Publishers, Dordrecht, 1994.
- [4] S. Gottwald, Fuzzy Sets and Fuzzy Logic. The Foundations of Application – from a Mathematical Point of View, Vieweg+Teubner Verlag, Braunschweig/Wiesbaden, 1993.
- [5] E. P. Klement, R. Mesiar, E. Pap, Triangular Norms, Kluwer Academic Publishers, Dordrecht, 2000.
- [6] E. Trillas, C. Alsina, E. Renedo, On some schemes of reasoning in fuzzy logic, New Math. Nat. Comput 7 (3) (2011) 433–451.

- [7] N. R. Vemuri, Investigations of fuzzy implications satisfying generalized hypothetical syllogism, *Fuzzy Sets and Systems* 323 (2017) 117–137.
- [8] L. A. Zadeh, Outline of a new approach to the analysis of complex systems and decision processes, *IEEE Trans. on Syst. Man and Cyber.* 3 (1973) 28–44.