

# Robust Optimization for the Counterpart Open Capacitated Vehicle Routing Problem With Time Windows

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## ABSTRACT

This paper investigates the robust optimization for the counterpart open capacitated vehicle routing problem (denoted by RCOCVRP) with soft time windows on waste transport problems. The uncertainty of volume of waste and travel time are considered. Garbage collection vehicles in carrying out the process of collecting garbage sometimes there are delays such as congestion and engine damage. We investigate the robust counterpart open capacitated vehicle routing problem with soft time windows models, where the waste transport vehicle does not need to return to the location of origin, the volume of waste carried by the vehicle does not exceed the vehicle capacity and consider the time windows. The robust counterpart open capacitated vehicle routing problem with soft time windows models solved by cheapest insertion heuristic algorithm. The results of the completion of the RCOCVRP model with soft time windows were completed with the GAMS software. Finally, the results of the calculation of the RCOCVRP model with soft time windows using the cheapest insertion heuristics algorithm are also compared with the GAMS software.

**Keywords:** Waste Transportation, Soft Time Windows, Robust Optimization, Travel Time.

## 1. INTRODUCTION

The problem of route of waste transportation vehicles is a major concern in the city that affects environmental problems. The problem of garbage transportation vehicle routes is a major concern in big cities which has an impact on environmental problems. The existence of residential areas, traditional markets, schools and offices which results in an increase in the volume of waste, such as in Kalidoni sub-district, Palembang city. The robust counterpart model of open capacitated vehicle routing problem (denoted by RCOCVRP) with soft time windows with soft time windows is a robust optimization for data that is uncertain [1], [2]. This model is applied to optimize the route of garbage collection vehicles with uncertain waste volume, time windows and travel time. Garbage collection vehicles depart from a node and do not return at that node. This differs from the classic vehicle routing problem (VRP). Classic VRP departs from an initial node and returns to initial node [3], [4]. Capacitated routing problem (CVRP) is also known as classic VRP related to vehicle capacity [5]–[7]. The problem of route

for transporting waste can be considered as a optimization problem [8]–[10]. Garbage collection vehicles carry garbage in several Temporary Disposal Sites (TDS) and end up at Final Disposal Sites (FDS). The process of transporting waste at some TDS sometimes occurs delays caused by congestion and machine damage. Garbage transport vehicles carry out the process of transporting waste from 06.00–10.00 and 16.00–20.00. The concept of this model is robust capacitated vehicle routing problem [11], robust open vehicle routing problem [12], robust vehicle routing problem with time windows [13]–[15] and robust counterpart open capacitated vehicle routing problem [16], [17]. The problem of route for transporting waste with uncertainty of waste volume can be considered as a robust optimization problem [17], [18], [2].

This study aims to optimize the route of the garbage transporting vehicle in the RCOCVRP model with soft time windows, limited by vehicle capacity, time windows and travel time. Time windows are the time intervals when the vehicle departs and arrives [9]. The RCOCVRP model with soft time windows was solved

using the cheapest insertion heuristic algorithm and GAMS software.

The cheapest insertion heuristic algorithm was first introduced to solve the Traveling Salesman Problem (TSP) and is known as the heuristic method [19]. This algorithm has been extended to VRP problems [20]. In this study, the cheapest insertion heuristic algorithm is used to optimize the route of the waste transport vehicle, including minimizing the distance and travel time. The concept of this algorithm is to insert each point in an arc from the starting point to the destination point [21]. The results of the cheapest insertion heuristic algorithm were solved by the GAMS software. GAMS (General Algebraic Modeling Software) is a very powerful optimization tool used to model, solve and analyze optimization problems [22]. Das and Kumar [23] uses GAMS to solve optimization problems and nonlinear programming (NLP) is used for solvers in GAMS such as CONOPT. Yuliza et al [1] has discussed the RCOCVRP model with soft time windows which was solved using LINGO and GAMS software. The resulting solution using GAMS is more optimal than using LINGO, indicated by a smaller calculation of the distance and travel time. In this study, linear programming (LP) was used for solvers in GAMS such as CPLEX.

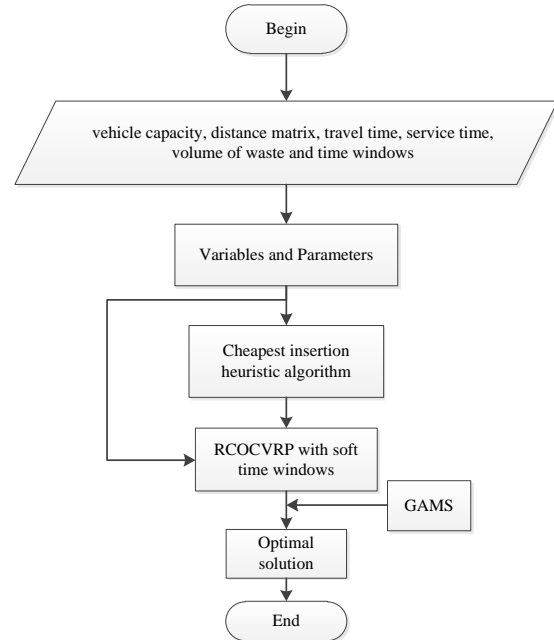
**2. METHOD**

Let  $G = (V, E)$  be a directed edge weighted graph, where  $V = \{1, 2, 3, \dots, n, n + 1\}$  denotes a set of vertices with nodes  $n + 1$  denotes the objective and  $E = \{(i, j) | i, j \in V, i \neq j\}$ . Non negative  $c_{ij}$  is represent the distance of arc  $(i, j) \in E$ . In this case the distance matrix is symmetric, in others word  $c_{ij} = c_{ji}$  for all  $i, j \in V, i \neq j$ . The cheapest insertion heuristic algorithm is as follows:

1. Route starts from the node of origin associated with the destination node.
2. A sub tour is generated between two vertex. A sub tour means a journey from a starting point and ending at a destination point.
3. Change one of the arc between two vertex with a combination of two arcs. Insertion  $k$  vertex on arc  $(i, j)$  with  $k$  vertex is not a member of sub tour and choose the smallest distance. Formula of the smallest distance could be:  
 $c_{ik} + c_{kj} - c_{ij}$  where  $c_{ik}$  is the distance from vertex  $i$  to  $k$ ,  $c_{kj}$  is the distance from vertex  $k$  to  $j$  and  $c_{ij}$  is the distance from vertex  $i$  to  $j$ .
4. Repeat step 3 until all the vertex are included in the sub tour.

In the case of the route of the garbage transporting vehicle, the vehicle departs from the starting point and does not return to that point so that the cheapest insertion heuristic algorithm is modified with limitations

on vehicle capacity and travel time related to time windows and service time. Service time is the time used to load and unload waste from the vehicle. A sub tour obtained from the cheapest insertion heuristic algorithm is solved using the CPLEX solver from GAMS. The RCOCVRP model with soft time windows was solved by the GAMS software. This research framework is as shown in Figure 1.



**Figure 1.** Research Framework

**3. RESULT AND DISCUSSION**

The research data was obtained from the Palembang City Environment Agency (PCEA) in the form of garbage transportation route data in the city of Palembang in September 2019. The data used in this study is the waste transportation route in Kalidoni sub-district, Palembang City. The following are some of the things that are limited:

1. It is assumed that the dump truck is 8 tons.
2. Assume the service time is approximately 10 minutes.
3. Assuming the time interval is 4 hours.

**Table 1.** Distance matrix (km)

Nodes	1	2	3	4	5	6	7	8
1	0	0.85	2.3	2.3	2.7	3.5	9.5	7.2
2		0	1.6	1.6	2	2.8	8,7	7
3			0	1.8	2.2	3	9	7.3
4				0	0.45	1.2	7.2	8.5
5					0	0.9	6.9	8.2
6						0	8.6	8.5
7							0	15
8								0

Kalidoni sub-district has 4 working areas, each working area has several TDS. Distance between each TDS and TDS and FDS as shown in Table 1. In this study, two examples of waste volume in working area 3 in each TDS are given such as Table 2 and time windows for each vertex as shown in Table 3.

**Table 2.** Volume of waste in the working area 3 (ton)

	1	2	3	4	5	6	7
Volume of waste_1	0.95	0.56	1.36	1.4	0.88	0.85	2
Volume of waste_2	2.3	1.2	2.4	2.4	1.5	2.6	1.3

**Table 3.** Time windows for working area 3

Vertex	Time windows
1	[07:00, 07:30]
2	[07:30, 08:00]
3	[08:00, 08:30]
4	[08:30, 09:00]
5	[09:00, 09:30]
6	[09:30, 10:00]
7	[10.00, 10.30]
8	[10.30, 11.00]

Variables and parameters in the working area 3 for  $i = \{1,2,3, \dots,8\}$  dan  $j = \{1,2,3, \dots,8\}$  as in Table 4 and 5. The optimal route RCOCVRP with soft time windows model on working area 3 uses the cheapest insertion heuristic algorithm as shown in Table 6 and Table 7 based on the volume of waste in Table 2.

**Table 4.** Variables and descriptions

Variable	Description
$x_{ij}$	Travelling from node $i$ to node $j$
$t_i$	Time arrived of the vehicle when serving node to $i$
$y_i$	Vehicle load when leaving node to $i$

**Table 5.** Parameters and descriptions

Parameter	Description
$c_{ij}$	distance from node $i$ to node $j$
$Q$	vehicle capacity
$q_i$	quantity of garbage transported at node to $i$

$t_{ij}$	travel time from node to $i$ to node to $j$
$a_j$	vehicle departure time from node to $j$
$b_i$	vehicle arrival time at node $i$

**Table 6.** The results of calculations using the cheapest insertion heuristic algorithm based on the volume of waste\_1

Garbage transport routes	Total distance (km)	Total travel time (minute)
TDS1 – TDS2 – TDS5 TDS6 – TDS7 – TDS4 – TDS3 – FDS	28.65	52.995

**Table 7.** The results of calculations using the cheapest insertion heuristic algorithm based on the volume of waste\_2.

Garbage transport routes	Total distance (km)	Total travel time (minute)
TDS1 – TDS2 – TDS5 – TDS3 – FDS	12.35	58.605
TDS4 – TDS6 – TDS7 – FDS	24.8	47.22

The model of RCOCVRP with soft time windows in work area 3 is based on data from Table 1, waste volume\_1 from Table 2 and time windows from Table 3 as follows:

Objective function

Minimize

$$0.85x_{12} + 2.3x_{13} + 2.3x_{14} + 2.7x_{15} + 3.5x_{16} + 9.5x_{17} + 7.2x_{18} + 0.85x_{21} + 1.6x_{23} + 1.6x_{24} + 2x_{25} + 2.8x_{26} + 8.7x_{27} + 7x_{28} + 2.3x_{31} + 1.6x_{32} + 1.8x_{34} + 2.2x_{35} + 3x_{36} + 9x_{37} + 7.3x_{38} + 2.3x_{41} + 1.6x_{42} + 1.8x_{43} + 0.45x_{45} + 1.2x_{46} + 7.2x_{47} + 8.5x_{48} + 2.7x_{51} + 2x_{52} + 2.2x_{53} + 0.45x_{54} + 0.9x_{56} + 6.9x_{57} + 8.2x_{58} + 3.5x_{61} + 2.8x_{62} + 3x_{63} + 1.2x_{64} + 0.9x_{65} + 8.6x_{67} + 8.5x_{68} + 9.5x_{71} + 8.7x_{72} + 9x_{73} + 7.2x_{74} + 6.9x_{75} + 8.6x_{76} + 15x_{78} + 7.2x_{81} + x_{782} + 7.3x_{83} + 8.5x_{84} + 8.2x_{85} + 8.5x_{86} + 15x_{87}$$

subject to

$$x_{12} + x_{13} + x_{14} + x_{15} + x_{16} + x_{17} + x_{18} = 1$$

$$x_{21} + x_{23} + x_{24} + x_{25} + x_{26} + x_{27} + x_{28} = 1$$

$$x_{31} + x_{32} + x_{34} + x_{35} + x_{36} + x_{37} + x_{38} = 1$$

$$\begin{aligned}
 x_{41} + x_{42} + x_{43} + x_{45} + x_{46} + x_{47} + x_{48} &= 1 \\
 x_{51} + x_{52} + x_{53} + x_{54} + x_{56} + x_{57} + x_{58} &= 1 \\
 x_{61} + x_{62} + x_{63} + x_{64} + x_{65} + x_{67} + x_{68} &= 1 \\
 x_{71} + x_{72} + x_{73} + x_{74} + x_{75} + x_{76} + x_{78} &= 1 \\
 x_{81} + x_{82} + x_{83} + x_{84} + x_{85} + x_{86} + x_{87} &= 1 \\
 x_{21} + x_{31} + x_{41} + x_{51} + x_{61} + x_{71} + x_{81} \\
 &\quad - x_{12} - x_{13} - x_{14} - x_{15} \\
 &\quad - x_{16} - x_{17} - x_{18} &= -1 \\
 x_{12} + x_{32} + x_{42} + x_{52} + x_{62} + x_{72} + x_{82} \\
 &\quad - x_{21} - x_{23} - x_{24} - x_{25} \\
 &\quad - x_{26} - x_{27} - x_{28} &= 0 \\
 x_{13} + x_{23} + x_{43} + x_{53} + x_{63} + x_{73} + x_{83} \\
 &\quad - x_{31} - x_{32} - x_{34} - x_{35} \\
 &\quad - x_{36} - x_{37} - x_{38} &= 0 \\
 x_{14} + x_{24} + x_{34} + x_{54} + x_{64} + x_{74} + x_{84} \\
 &\quad - x_{41} - x_{42} - x_{43} - x_{45} \\
 &\quad - x_{46} - x_{47} - x_{48} &= 0 \\
 x_{15} + x_{25} + x_{35} + x_{45} + x_{65} + x_{75} + x_{85} \\
 &\quad - x_{51} - x_{52} - x_{53} - x_{54} \\
 &\quad - x_{56} - x_{57} - x_{58} &= 0 \\
 x_{16} + x_{26} + x_{36} + x_{46} + x_{56} + x_{76} + x_{86} \\
 &\quad - x_{61} - x_{62} - x_{63} - x_{64} \\
 &\quad - x_{65} - x_{67} - x_{68} &= 0 \\
 x_{17} + x_{27} + x_{37} + x_{47} + x_{57} + x_{67} + x_{68} \\
 &\quad - x_{71} - x_{72} - x_{73} - x_{74} \\
 &\quad - x_{75} - x_{76} - x_{78} &= 0 \\
 x_{18} + x_{28} + x_{38} + x_{48} + x_{58} + x_{68} + x_{78} \\
 &\quad - x_{81} - x_{82} - x_{83} - x_{84} \\
 &\quad - x_{85} - x_{86} - x_{87} &= 1
 \end{aligned}$$

$$950 \leq y_1 \leq 8000$$

$$560 \leq y_2 \leq 8000$$

$$1360 \leq y_3 \leq 8000$$

$$1400 \leq y_4 \leq 8000$$

$$880 \leq y_5 \leq 8000$$

$$850 \leq y_6 \leq 8000$$

$$2000 \leq y_7 \leq 8000$$

$$0 \leq y_8 \leq 8000$$

$$y_1 - y_2 + 8000x_{12} \leq 7050$$

$$y_2 - y_3 + 8000x_{23} \leq 7740$$

$$y_3 - y_4 + 8000x_{34} \leq 6640$$

$$y_4 - y_5 + 8000x_{45} \leq 6600$$

$$y_5 - y_6 + 8000x_{56} \leq 7120$$

$$y_6 - y_7 + 8000x_{67} \leq 7150$$

$$y_7 - y_8 + 8000x_{78} \leq 8000$$

$$7 \leq t_1 \leq 7,5$$

$$7,5 \leq t_2 \leq 8$$

$$8 \leq t_3 \leq 8,5$$

$$8,5 \leq t_4 \leq 9$$

$$9 \leq t_5 \leq 9,5$$

$$9,5 \leq t_6 \leq 10$$

$$10 \leq t_7 \leq 10,5$$

$$10,5 \leq t_8 \leq 11$$

$$t_1 - t_2 + 0,02125x_{12} + 0,167 \leq 0$$

$$t_2 - t_3 + 0,04x_{23} + 0,167 \leq 0$$

$$t_3 - t_4 - 0,0225x_{34} + 0,167 \leq -0,3$$

$$t_4 - t_5 + 0,01125x_{45} + 0,167 \leq 0$$

$$t_5 - t_6 - 0,2725x_{56} + 0,167 \leq -0,3$$

$$t_6 - t_7 + 0,215x_{67} + 0,167 \leq 0$$

$$t_7 - t_8 - 0,2625x_{78} + 0,167 \leq -0,3$$

## 4. CONCLUSION

The RCOCVRP model with a soft time window with uncertainty of waste volume and travel time is solved using the cheapest insertion heuristic algorithm which yields various solutions. The RCOCVRP model with a soft time window with uncertain waste volume and trip times completed using the CPLEX solver from GAMS yields the optimal solution. Although the total distance and the total travel time calculated using the cheapest insertion heuristic algorithm are smaller than the results from GAMS.

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