

Closure Systems as a Fuzzy Extension of Meet-subsemilattices

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Abstract

An extension of the notion of closure system is studied adapting the idea of meet-subsemilattice to a complete fuzzy lattice. Results relating closure operators and closure systems in the classical case are extended properly to this framework. This definition is proved to be equivalent to the most used definition given by Bělohlávek on the fuzzy powerset lattice. A definition of fuzzy closure system is presented and related to closure systems.

1 Introduction

Closure operators have an important role in pure, applied mathematics and computer science [10]. Fuzzy closure operators [1, 4] appear in several areas of fuzzy logic and its applications, including fuzzy mathematical morphology [13, 18], fuzzy relational equations [12], approximate reasoning [6, 11] and fuzzy logic in narrow sense [15], and its applications such as fuzzy logic programming [17] or formal concept analysis of data with fuzzy attributes [20].

As in the classical case, there is a bijection which relates the notions of closure operator and closure system. We can find distinct definitions of closure system depending on the ordered structure on which the fuzzy closure operator is defined. As a consequence, the notion of fuzzy closure system has been defined on L -ordered sets [14], on fuzzy preposets [8] and fuzzy preordered structures [9].

This paper is a continuation on the study of fuzzy closure systems done in [19], where the underlying structure was a Heyting algebra. Our aim is to study the concept of closure system on complete fuzzy lattices. One can find several extensions of the notion of lattice

and complete lattice to the fuzzy framework. Basically, these extensions go in two main ways.

On the one hand, some authors define a fuzzy lattice as a fuzzy subset of a (classical) lattice L that is compatible with the structure of the lattice in some manner [21, 23]. This particular extension could be called that of *fuzzy sublattices*. Closure systems on fuzzy sublattices were introduced in [22].

On the other hand, sometimes a fuzzy lattice is defined as a (crisp) set endowed with a fuzzy order relation satisfying certain conditions related to the existence of generalized infimum and supremum. The concept of *complete fuzzy lattice* considered in this work belongs to this class. It was introduced in [5] with the name of *completely lattice L -ordered set*, although it has been called other names, for instance, *fuzzy complete lattice* in [16]. A particular case of fuzzy complete lattice is the L -powerset with the so-called subthood degree relation S . Closure systems on L -powerset lattices were introduced by Bělohlávek in [1].

The starting point of this paper is the definition of closure system, in the framework of complete fuzzy lattices, provided in [3]. This extension is inspired by the notion of closure system in the case of classical complete lattices, where they are \wedge -subsemilattices. As awaited, the relationship between fuzzy closure operators and closure systems is studied. After a first section of preliminaries, the second section includes the definition of closure system in a fuzzy framework, which extends the classical closure system property of being a \wedge -subsemilattice. A one-to-one relation between closure systems and fuzzy closure operators is proved, i.e., closure systems induce closure operators and closure operators induce closure systems in a desirable manner. In addition, a characterization of closure system in terms of the minimum of certain sets is proved. The core of the next section is the study of the relationship between the definition of closure system and the one of L -closure system given by Bělohlávek

in [1]. The main result within this section is the equivalence of both definitions on L -powerset lattices. Furthermore, the notion of fuzzy closure system in this framework, which is not as straightforward as the one on Heyting algebras provided in [19], is introduced and its good properties are shown. In the last section, the conclusions of our research are shown and some hints of future work are presented.

2 Preliminaries

The framework to which we are going to generalize the notion of closure system is introduced in this section. It has been chosen with the idea of being as general as possible and thus have a wider range of possible applications. Specifically, we recall the notions of complete residuated lattices [2, 15], fuzzy poset and fuzzy complete lattice, and some basics results that will be needed throughout the paper [5, 16].

Throughout this paper, let $\mathbb{L} = (L, \wedge, \vee, \otimes, \rightarrow, 0, 1)$ be a complete residuated lattice, which is an algebra where

- $(L, \wedge, \vee, 0, 1)$ is a complete lattice with 0 and 1 being the least and the greatest elements of L , respectively,
- $(L, \otimes, 1)$ is a commutative monoid (i.e., \otimes is commutative, associative, and 1 is neutral with respect to \otimes), and
- \otimes and \rightarrow satisfy the so-called *adjointness property*: for all $a, b, c \in L$, we have that $a \otimes b \leq c$ iff $a \leq b \rightarrow c$.

This structure is utilized in mathematical fuzzy logics and their applications as structures of truth degrees with \otimes and \rightarrow used as truth functions of *fuzzy conjunction* and *fuzzy implication*, respectively [15]. The unit interval with the Łukasiewicz, Gödel and Goguen pairs of t-norms and implications are examples of residuated complete lattices.

An \mathbb{L} -set is a mapping $X: U \rightarrow L$ from the universe set U to the membership values set L , where $X(u)$ means the degree in which u belongs to X . The set of \mathbb{L} -sets on the universe U is denoted by L^U . A crisp set is considered to be a particular case of \mathbb{L} -set by using its characteristic mapping $X: U \rightarrow \{0, 1\}$ with $X(u) = 1$ iff $u \in X$.

Operations with \mathbb{L} -sets are defined element-wise. For instance, $A \otimes B \in L^U$ is defined as $(A \otimes B)(u) = A(u) \otimes B(u)$ for all $u \in U$. The so-called *subsethood degree*

relation is defined as $S: L^U \times L^U \rightarrow L$ where

$$S(A, B) = \bigwedge_{x \in U} (A(x) \rightarrow B(x)).$$

Obviously, $S(A, B) = 1$ if and only if $A(x) \leq B(x)$ for all $x \in U$. In this case, A is said to be a fuzzy subset of B and is denoted by $A \subseteq B$.

Binary \mathbb{L} -relations (binary fuzzy relations) on a set U can be thought of as \mathbb{L} -sets on the universe $U \times U$. That is, a binary \mathbb{L} -relation on U is a mapping $\rho \in L^{U \times U}$ assigning to each $x, y \in U$ a truth degree $\rho(x, y) \in L$ (a degree to which x and y are related by ρ).

For ρ being a binary \mathbb{L} -relation in U , we say that

- ρ is *reflexive* if $\rho(x, x) = 1$ for all $x \in U$.
- ρ is *symmetric* if $\rho(x, y) = \rho(y, x)$ for all $x, y \in U$.
- ρ is *antisymmetric* if $\rho(x, y) \otimes \rho(y, x) = 1$ implies $x = y$ for all $x, y \in U$.
- ρ is *transitive* if $\rho(x, y) \otimes \rho(y, z) \leq \rho(x, z)$ for all $x, y, z \in U$.

Definition 1. Given a non-empty set A and a binary \mathbb{L} -relation ρ on A , the pair $\mathbb{A} = (A, \rho)$ is said to be a

- *fuzzy preposet* if ρ is a *fuzzy preorder*, i.e. if ρ is reflexive and transitive;
- *fuzzy poset* if ρ is a *fuzzy order*, i.e. if ρ is reflexive, antisymmetric and transitive.

A typical example of fuzzy poset is (L^U, S) .

As in the crisp case, any order \mathbb{L} -relation implicitly defines an equivalence \mathbb{L} -relation that is called *symmetric kernel relation*. In the fuzzy case, this equivalence \mathbb{L} -relation usually replaces the notion of equality in the fuzzy poset.

Definition 2. Given a fuzzy preposet $\mathbb{A} = (A, \rho)$, the *symmetric kernel relation* is defined as $\approx: A \times A \rightarrow L$ where $(a \approx b) = \rho(a, b) \wedge \rho(b, a)$ for all $a, b \in A$.

Proposition 3. Given a fuzzy preposet $\mathbb{A} = (A, \rho)$, the *symmetric kernel relation* \approx is a *fuzzy equivalence relation*, that is, it is a reflexive, symmetric and transitive fuzzy relation.

A usual way to define fuzzy algebras is to consider as an underlying structure a pair which consists of a set and a tolerance or equivalence relation on it. Thus, an alternative definition of fuzzy poset that can be found in the literature [9] is given by a tuple (A, \approx, ρ) where \approx is a fuzzy equivalence relation on A and ρ is a fuzzy order that is compatible with \approx . In [24], it is shown that both definitions of fuzzy poset are equivalent.

To present the notion of fuzzy lattice we need to generalize those of upper (lower) bound and supremum (infimum).

Definition 4. Given a fuzzy preposet $\mathbb{A} = (A, \rho)$ and a fuzzy set $X \in L^A$, we define the *down-cone* of X and the *up-cone* of X , respectively, as the fuzzy sets $X_\rho, X^\rho \in L^A$ where, for all $a \in A$,

$$X_\rho(a) = \bigwedge_{x \in A} (X(x) \rightarrow \rho(a, x)) \quad \text{and}$$

$$X^\rho(a) = \bigwedge_{x \in A} (X(x) \rightarrow \rho(x, a)).$$

Thus, $X_\rho(a)$ and $X^\rho(a)$ can be seen as the degree to which a is an upper bound and lower bound of X , respectively.

Notice that for singletons, $X = \{x\}$, we will omit the brackets for simplicity of the notation, thus $x^\rho(a) = \rho(x, a)$ and $x_\rho(a) = \rho(a, x)$.

Definition 5. Let $\mathbb{A} = (A, \rho)$ be a fuzzy preposet and a fuzzy subset X of A .

An element $m \in A$ is a *p-minimum* of X iff

- i. $X(m) = 1$, and
- ii. $X \subseteq m^\rho$, i.e., $X(x) \leq \rho(m, x)$, for all $x \in A$.

An element $M \in A$ is a *p-maximum* of X iff

- i. $X(M) = 1$, and
- ii. $X \subseteq M_\rho$, i.e., $X(x) \leq \rho(x, M)$, for all $x \in A$.

The crisp set of p-minima (resp. p-maxima) for X will be denoted $\text{p-min}(X)$ (resp. $\text{p-max}(X)$). Notice that, due to the absence of antisymmetry, this set could have more than one element. If \mathbb{A} is a poset, then every non-empty set of p-minima and p-maxima is a singleton whose element will be called the minimum (resp. maximum) of X , denoted by $\min(X)$ (resp. $\max(X)$).

Definition 6. Let $\mathbb{A} = (A, \rho)$ be a fuzzy preposet and $X \in L^A$. An element $a \in A$ is said to be *infimum* (resp. *supremum*) of X if the following conditions hold:

- i. $X_\rho(a) = 1$ (resp. $X^\rho(a) = 1$).
- ii. $X_\rho(x) \leq \rho(x, a)$ (resp. $X^\rho(x) \leq \rho(a, x)$), for all $x \in A$.

The following propositions, which are straightforward, provide useful characterizations of infimum, supremum, p-minimum and p-maximum.

Proposition 7. Let $\mathbb{A} = (A, \rho)$ be a fuzzy preposet and $X \in L^A$. An element $a \in A$ is infimum (resp. supremum) of X if and only if, for all $x \in A$,

$$\rho(x, a) = X_\rho(x) \quad (\text{resp. } \rho(a, x) = X^\rho(x)).$$

Proposition 8. Let $\mathbb{A} = (A, \rho)$ be a fuzzy preposet and $X \in L^A$. An element $m \in A$ is a p-minimum (resp. p-maximum) of X if and only if m is an infimum (resp. a supremum) of X and $X(m) = 1$.

Notice that the supremum (resp. infimum) needs not be unique in an arbitrary preposet, but if \mathbb{A} is a poset and due to antisymmetry, it is not difficult to see that, if a supremum (resp. infimum) of X exists, it is unique. In such case, we will denote it by $\bigsqcup X$ (resp. $\bigsqcap X$).

Definition 9 ([5]). We say that a fuzzy poset (A, ρ) is a complete fuzzy lattice if every fuzzy subset $X \in L^A$ has supremum and infimum.

The pair (L^U, S) is an example of complete fuzzy lattice, which is called the *L-powerset lattice of U*. This fact follows easily from [2, Theorem 5.63].

Notice that, if (A, ρ) is a complete fuzzy lattice, then (A, \leq) , with $a \leq b$ iff $\rho(a, b) = 1$, is a complete lattice. Therefore, there exist elements that are minimum and maximum, which we denote by \perp and \top respectively.

We conclude this section with the usual definition of closure operator on a fuzzy preposet.

Definition 10. Given a fuzzy preposet $\mathbb{A} = (A, \rho)$, a mapping $c: A \rightarrow A$ is said to be a *closure operator* on \mathbb{A} if the following conditions hold:

- i. $\rho(a, b) \leq \rho(c(a), c(b))$, for all $a, b \in A$
- ii. $\rho(a, c(a)) = 1$, for all $a \in A$
- iii. $\rho(c(c(a)), c(a)) = 1$, for all $a \in A$.

Conditions i and ii are well-known and are called isotony and inflationarity, respectively. Observe that condition iii could be replaced by $(c(c(a)) \approx c(a)) = 1$, and, thus, if \mathbb{A} is a fuzzy poset, a closure operator is idempotent in a classical sense, i.e., $c(c(a)) = c(a)$ for all $a \in A$.

3 Closure systems on fuzzy lattices

A generalization of the classical definition of closure system on complete lattices was proposed in [3], where they are considered as sets with the \wedge -subsemilattice structure.

Definition 11. Let (A, ρ) be a fuzzy complete lattice. A crisp subset $\mathcal{F} \subseteq A$ is said to be a closure system if $\bigsqcap X \in \mathcal{F}$ for any fuzzy subset X of \mathcal{F} .

Notice that the empty set is contained in every set and $\bigcap \emptyset = \top$. Therefore, \top is in every closure system.

Example 1. Let $\mathbb{L} = (\{0, 0.5, 1\}, \wedge, \vee, \otimes, \rightarrow, 0, 1)$ be the three-valued Łukasiewicz residuated lattice, and (A, ρ) be the fuzzy lattice where $A = \{\perp, a, b, c, d, e, \top\}$ and $\rho: A \times A \rightarrow L$ is the fuzzy relation order described by the following table:

ρ	\perp	a	b	c	d	e	\top
\perp	1	1	1	1	1	1	1
a	0.5	1	0.5	1	1	1	1
b	0.5	0.5	1	1	1	1	1
c	0.5	0.5	0.5	1	1	1	1
d	0	0.5	0	0.5	1	0.5	1
e	0	0	0.5	0.5	0.5	1	1
\top	0	0	0	0.5	0.5	0.5	1

The set $\mathcal{F} = \{c, \top\}$ is a closure system because, for any fuzzy subset X of \mathcal{F} , we have that $\bigcap X = c$ when $X(c) = 1$, and $\bigcap X = \top$ otherwise.

Now we provide a proof for the following theorem that was stated in [3]. It shows the existence of a bijection between closure systems and closure operators, as in the classical case.

Theorem 12. Let $\mathbb{A} = (A, \rho)$ be a complete fuzzy lattice.

- i. If \mathcal{F} is a closure system on \mathbb{A} , then the mapping $c_{\mathcal{F}}: A \rightarrow A$ defined as $c_{\mathcal{F}}(x) = \bigcap (x^{\rho} \cap \mathcal{F})$ is a closure operator on \mathbb{A} .
- ii. If $c: A \rightarrow A$ is a closure operator on \mathbb{A} , then $\mathcal{F}_c = \{x \in A \mid c(x) = x\}$ is a closure system on \mathbb{A} .
- iii. If \mathcal{F} is a closure system on \mathbb{A} , then $\mathcal{F}_{c_{\mathcal{F}}} = \mathcal{F}$.
- iv. If $c: A \rightarrow A$ is a closure operator on \mathbb{A} , then $c_{\mathcal{F}_c} = c$.

Proof. First, assume that \mathcal{F} is a closure system and, therefore, $c_{\mathcal{F}}(x) \in \mathcal{F}$ for all $x \in A$.

$$\begin{aligned} \rho(x, c_{\mathcal{F}}(x)) &= \rho\left(x, \bigcap (x^{\rho} \cap \mathcal{F})\right) = \\ &= (x^{\rho} \cap \mathcal{F})_{\rho}(x) = \bigwedge_{a \in A} [(x^{\rho} \cap \mathcal{F})(a) \rightarrow \rho(x, a)] \end{aligned}$$

Since $(x^{\rho} \cap \mathcal{F})(a)$ is $\rho(x, a)$ if $a \in \mathcal{F}$ and 0 otherwise, we get,

$$\begin{aligned} &\bigwedge_{a \in A} [(x^{\rho} \cap \mathcal{F})(a) \rightarrow \rho(x, a)] = \\ &= \bigwedge_{a \in \mathcal{F}} [\rho(x, a) \rightarrow \rho(x, a)] = 1, \text{ for all } x \in A \end{aligned}$$

Hence, $\rho(x, c_{\mathcal{F}}(x)) = 1$, for all $x \in A$, i.e. $c_{\mathcal{F}}$ is inflationary.

By the properties of infima, for all $a \in A$ and $X \in L^A$, we have $X(a) \leq \rho(\bigcap X, a)$. In particular, for $X = x^{\rho} \cap \mathcal{F}$, we have that

$$\rho(x, v) = (x^{\rho} \cap \mathcal{F})(v) \leq \rho(c_{\mathcal{F}}(x), v), \quad (1)$$

for all $x \in A, v \in \mathcal{F}$.

Now, for any $x, y \in A$, since $c_{\mathcal{F}}$ is inflationary and $c_{\mathcal{F}}(y) \in \mathcal{F}$, it follows,

$$\begin{aligned} \rho(x, y) &= \rho(x, y) \otimes \rho(y, c_{\mathcal{F}}(y)) \leq \\ &\leq \rho(x, c_{\mathcal{F}}(y)) \leq \rho(c_{\mathcal{F}}(x), c_{\mathcal{F}}(y)). \end{aligned}$$

Therefore, $c_{\mathcal{F}}$ is isotone.

Let $x \in A$. Since, by hypothesis, $c_{\mathcal{F}}(x) \in \mathcal{F}$, and by using (1), we get

$$1 = \rho(c_{\mathcal{F}}(x), c_{\mathcal{F}}(x)) \leq \rho(c_{\mathcal{F}}(c_{\mathcal{F}}(x)), c_{\mathcal{F}}(x)).$$

Therefore, $\rho(c_{\mathcal{F}}(c_{\mathcal{F}}(x)), c_{\mathcal{F}}(x)) = 1$, for all $x \in A$, i.e. $c_{\mathcal{F}}$ is idempotent; and we have proved that it is a closure operator.

For ii, assume c is a closure operator in (A, ρ) and let $\mathcal{F}_c = \{x \in A \mid c(x) = x\}$. For every fuzzy subset $X \subseteq \mathcal{F}_c$, we denote $m = \bigcap X$ and prove $m \in \mathcal{F}_c$.

By Proposition 7, we have that

$$\begin{aligned} \rho(c(m), m) &= X_{\rho}(c(m)) = \\ &= \bigwedge_{a \in A} (X(a) \rightarrow \rho(c(m), a)) \end{aligned}$$

Now, since $X \subseteq \mathcal{F}_c$, we have that $X(a) = 0$ for all $a \notin \mathcal{F}_c$ and, by using it and the fact that $c(s) = s$ for all $s \in \mathcal{F}_c$, we obtain

$$\begin{aligned} \rho(c(m), m) &= \bigwedge_{s \in \mathcal{F}_c} (X(s) \rightarrow \rho(c(m), s)) = \\ &= \bigwedge_{s \in \mathcal{F}_c} (X(s) \rightarrow \rho(c(m), c(s))) \end{aligned}$$

Therefore, $\rho(c(m), m) = 1$ because, since m is the infimum of X and c is isotone, we have $X(s) \leq \rho(m, s) \leq \rho(c(m), c(s))$. Finally, since c is inflationary, and by antisymmetry, we get $m = c(m) \in \mathcal{F}_c$.

Item iii is proved by double inclusion. On the one hand, if $s \in \mathcal{F}$, then by the first property of the infimum

$$1 = \rho(s, s) = (s^{\rho} \cap \mathcal{F})(s) \leq \rho(c_{\mathcal{F}}(s), s).$$

Also, by inflationarity of the closure operator we had $\rho(s, c_{\mathcal{F}}(s)) = 1$. Hence, using antisymmetry $\rho(s, c_{\mathcal{F}}(s)) \otimes \rho(c_{\mathcal{F}}(s), s) = 1$ implies $s = c_{\mathcal{F}}(s)$. Then, $s \in \mathcal{F}_{c_{\mathcal{F}}}$ and we have proved $\mathcal{F} \subseteq \mathcal{F}_{c_{\mathcal{F}}}$.

On the other hand, if $s \in \mathcal{F}_{c_{\mathcal{F}}}$, this means $\rho(c_{\mathcal{F}}(s), s) = 1$. By inflationarity we have

$\rho(s, c_{\mathcal{F}}(s)) = 1$. Hence by antisymmetry $s = c_{\mathcal{F}}(s)$ and by definition of $c_{\mathcal{F}}$, $s = c_{\mathcal{F}}(s) \in \mathcal{F}$. Thus, $\mathcal{F}_{c_{\mathcal{F}}} \subseteq \mathcal{F}$.

In order to prove item iv, i.e. $c(x) = c_{\mathcal{F}_c}(x)$, it suffices to show $c(x)$ is the infimum of $(x^p \cap \mathcal{F}_c)$.

First, for all $y \in A$, we have

$$\begin{aligned} (x^p \cap \mathcal{F}_c)_\rho(y) &= \bigwedge_{a \in A} [(x^p \cap \mathcal{F}_c)(a) \rightarrow \rho(y, a)] = \\ &= \bigwedge_{a \in \mathcal{F}_c} [\rho(x, a) \rightarrow \rho(y, a)] \\ &\leq \rho(x, c(x)) \rightarrow \rho(y, c(x)) = \rho(y, c(x)). \end{aligned}$$

Moreover, given $y \in A$ and $a \in \mathcal{F}_c$, we obtain

$$\rho(y, c(x)) \leq \rho(c(x), a) \rightarrow \rho(y, a) \stackrel{(1)}{\leq} \rho(x, a) \rightarrow \rho(y, a)$$

and, since it is true for all $a \in \mathcal{F}_c$, we have that

$$\rho(y, c(x)) \leq \bigwedge_{a \in \mathcal{F}_c} [\rho(x, a) \rightarrow \rho(y, a)] = (x^p \cap \mathcal{F}_c)_\rho(y).$$

Thus, by Proposition 7, $c(x) = \prod(x^p \cap \mathcal{F}_c)$, for all $x \in A$ and $c = c_{\mathcal{F}_c}$ is proved. \square

The previous theorem establishes a bijection between closure systems and closure operators in complete fuzzy lattices. Given a closure operator $c: A \rightarrow A$, the set $\mathcal{F}_c = \{x \in A \mid c(x) = x\}$ is the closure system associated to c , and, given a closure system \mathcal{F} , the mapping $c_{\mathcal{F}}: A \rightarrow A$ defined as $c_{\mathcal{F}}(x) = \prod(x^p \cap \mathcal{F})$ is the closure operator associated to \mathcal{F} . This last construction leads to a characterization of closure systems in terms of minima.

Proposition 13. *Let $\mathbb{A} = (A, \rho)$ be a complete fuzzy lattice. A crisp set $\mathcal{F} \subseteq A$ is a closure system if and only if the element $\min(x^p \cap \mathcal{F})$ exists for all $x \in A$.*

Proof. Assume \mathcal{F} is a closure system and let $x \in A$. Consider the set $x^p \cap \mathcal{F}$, which is a subset of \mathcal{F} . Since \mathcal{F} is a closure system, $\prod(x^p \cap \mathcal{F})$ is an element of \mathcal{F} . Furthermore, from Theorem 12, we have that $\rho(x, \prod(x^p \cap \mathcal{F})) = 1$ and $(x^p \cap \mathcal{F})(\prod(x^p \cap \mathcal{F})) = 1$, so it is the minimum of the set.

Conversely, assume the set \mathcal{F} is such that for all $x \in A$, $\min(x^p \cap \mathcal{F})$ exists.

Let $X \subseteq \mathcal{F}$ and consider the set $(\prod X)^p \cap \mathcal{F}$. Following the assumption, the element $m = \min((\prod X)^p \cap \mathcal{F})$ exists. We claim $m = \prod X$.

By definition of minimum, $((\prod X)^p \cap \mathcal{F})(m) = 1$, i.e., $\rho(\prod X, m) = 1$ and $m \in \mathcal{F}$.

By definition of infimum, $X(x) \leq \rho(\prod X, x)$, for all $x \in A$ and $X_\rho(x) \leq \rho(x, \prod X)$, for all $x \in A$.

$$X_\rho(m) = \bigwedge_{x \in A} X(x) \rightarrow \rho(m, x) \stackrel{X \subseteq \mathcal{F}}{\equiv} \bigwedge_{s \in \mathcal{F}} X(s) \rightarrow \rho(m, s).$$

Let $s \in \mathcal{F}$, then

$$X(s) \leq \rho(\prod X, s) = (\prod X^p \cap \mathcal{F})(s) \leq \rho(m, s)$$

Hence, $X_\rho(m) = 1 \leq \rho(m, \prod X)$, so $\prod X = m \in \mathcal{F}$. \square

4 Equivalence with L-closure systems in powerset lattices

The most cited definition of closure system in a fuzzy setting was given by Bělohlávek in [1], where the notions of L_K -closure operator and L_K -closure system on \mathbb{L} -ordered sets were introduced, K being a filter of the residuated lattice \mathbb{L} . Moreover, the product and residuum of \mathbb{L} have been extended to an external operation on the powerset lattice as follows:

$$\begin{aligned} (a \otimes A)(x) &= a \otimes A(x) \text{ and} \\ (a \rightarrow A)(x) &= a \rightarrow A(x). \end{aligned} \quad (2)$$

for all $a, x \in L, A \in L^U$.

Notice that these operations involve elements of L and elements of L^U , hence they are particularly defined on the powerset lattice.

Definition 14 ([1]). Let K be a filter in L . A system $\mathcal{F} = \{A_i \in L^X \mid i \in I\}$ is called closed under S_K -intersections iff for each $A \in L^U$ it holds that

$$\bigcap_{i \in I, S(A, A_i) \in K} (S(A, A_i) \rightarrow A_i) \in \mathcal{F}$$

A system closed under S_K -intersections will be called an L_K -closure system. For $K = L$ the subscript will be omitted.

Remark 1. In general, being closed under arbitrary intersections –i.e. $\bigcap_{i \in J} A_i \in \mathcal{F}$ for all $J \subseteq I$ – is a weaker condition than being closed under S_K -intersections. See [1, Remark 3.3 (2)].

Our aim in this paper has been to highlight the definition of closure system on complete fuzzy lattices, not specifically over an L -powerset. Accordingly, the use of the external operations given in (2) are circumvented. To avoid confusion, we will rename Bělohlávek's L -closure systems as S -closure systems.

The following lemma is a technical result which is used in the proof of the next theorem.

Lemma 15. Let (L^U, S) be a L -powerset lattice. For all $A, B \in L^U$ we have that $A \subseteq S(A, B) \rightarrow B$ i.e. $S(A, S(A, B) \rightarrow B) = 1$.

Proof. It is a consequence of $S(A, B) \otimes A(x) \leq B(x)$ for all $x \in U$. \square

Theorem 16. Let (L^U, S) be a L -powerset lattice and $\mathcal{F} \subseteq L^U$. Then, \mathcal{F} is a closure system if and only if \mathcal{F} is an S -closure system. In addition, if \mathcal{F} is a closure system then, for all $A \in L^U$,

$$\min(A^S \cap \mathcal{F}) = \bigcap_{C \in \mathcal{F}} (S(A, C) \rightarrow C). \quad (3)$$

Proof. By Proposition 13, it is sufficient to prove that, in both cases, if \mathcal{F} is a closure system or an L -closure system, the equality (3) holds.

On the one hand, assuming $\mathcal{F} \subseteq L^U$ is a closure system, let us prove (3). By Proposition 13 the element $M_A = \min(A^S \cap \mathcal{F})$ exists, for all $A \in L^U$.

Since M_A is a minimum it verifies $(A^S \cap \mathcal{F})(M_A) = 1$ and, for all $B \in L^U$, $(A^S \cap \mathcal{F})(B) \leq S(M_A, B)$. That is, $M_A \in \mathcal{F}$, $S(A, M_A) = 1$ and, for all $C \in \mathcal{F}$, we have $S(A, C) \leq S(M_A, C)$.

Since $M_A \in \mathcal{F}$ and $S(A, M_A) = 1$, we have that

$$\bigcap_{C \in \mathcal{F}} (S(A, C) \rightarrow C) \subseteq S(A, M_A) \rightarrow M_A = M_A.$$

Moreover, by Lemma 15, and since $S(A, C) \leq S(M_A, C)$ for all $C \in \mathcal{F}$, we have that

$$M_A \subseteq \bigcap_{C \in \mathcal{F}} (S(M_A, C) \rightarrow C) \subseteq \bigcap_{C \in \mathcal{F}} (S(A, C) \rightarrow C).$$

Thus, $\bigcap_{C \in \mathcal{F}} (S(A, C) \rightarrow C) = M_A \in \mathcal{F}$, for all $A \in L^U$ and \mathcal{F} is an S -closure system.

On the other hand, assume \mathcal{F} is an S -closure system, that is, $\bigcap_{C \in \mathcal{F}} (S(A, C) \rightarrow C) \in \mathcal{F}$ for all $A \in L^U$. Let $F_A = \bigcap_{C \in \mathcal{F}} (S(A, C) \rightarrow C)$ and prove $F_A = \min(A^S \cap \mathcal{F})$. First, $F_A \in \mathcal{F}$ and by Lemma 15, we have that

$$\begin{aligned} S(A, F_A) &= S\left(A, \bigcap_{C \in \mathcal{F}} (S(A, C) \rightarrow C)\right) = \\ &= \bigwedge_{C \in \mathcal{F}} S(A, S(A, C) \rightarrow C) = 1. \end{aligned}$$

Finally, for all $C \in \mathcal{F}$, we have that $F_A \subseteq S(A, C) \rightarrow C$ and, therefore,

$$S(A, C) \leq \bigwedge_{x \in U} (F_A(x) \rightarrow C(x)) = S(F_A, C).$$

Thus, for all $A \in L^U$, $F_A = \min(A^S \cap \mathcal{F})$, hence \mathcal{F} is a closure system. \square

Thus, in an L -powerset lattice (L^U, S) , a closure system is a crisp subset $\mathcal{F} \subseteq L^U$ which is closed under arbitrary infima of fuzzy subsets ($X \in L^{\mathcal{F}}$ implies $\bigcap X \in \mathcal{F}$). This implies that \mathcal{F} is closed under arbitrary crisp intersections ($X \subseteq \mathcal{F}$ implies $\bigcap X \in \mathcal{F}$). Nevertheless, as stated in Remark 1, the reciprocal implication does not hold. We illustrate it with the following example.

Example 2. Consider the universe set $U = \{a, b\}$, the lattice L be the three-valued Łukasiewicz structure and (L^U, S) the L -powerset lattice of U . Then, the set $\mathcal{F} = \{\{a/1, b/0.5\}, \{a/1, b/1\}\}$ is a closure system. A fuzzy subset of \mathcal{F} can be described as $X = \{\{a/1, b/0.5\}/x, \{a/1, b/1\}/y\}$ where $x, y \in L$.

The infimum $\bigcap X$ is $\{a/1, b/0.5\}$ if $x = 1$ and $\bigcap X = \{a/1, b/1\}$ otherwise. Hence, \mathcal{F} is a closure system.

As proved in Theorem 16, closure systems and S -closure systems are the same concept within L -powerset lattices. Let us show that \mathcal{F} in this example is indeed an S -closure system.

In this particular case, $\bigcap_{C \in \mathcal{F}} S(A, C) \rightarrow C \in \mathcal{F}$ since for all $A \in L^U$ it is either $\{a/1, b/0.5\}$ or $\{a/1, b/0.5\}$.

Notice that these properties are stronger than being closed for arbitrary intersections. Consider the set $\mathcal{H} = \{\{a/0, b/0\}, \{a/1, b/1\}\}$, this set is clearly closed under intersections but it does not satisfy the closure system nor the S -closure system definition.

Consider $A = \{a/0.5, b/0.5\}$, then

$$\bigcap_{C \in \mathcal{H}} S(A, C) \rightarrow C = A \notin \mathcal{H}.$$

Thus, \mathcal{H} is not an S -closure system. Moreover, \mathcal{H} is not a closure system either since for $B = \{\{a/0, b/0\}/0.5\} \subseteq \mathcal{H}$, we have $\bigcap B = \{a/0.5, b/0.5\} \notin \mathcal{H}$. Thus, \mathcal{H} is not a closure system. The calculations of the infima and the intersections are left as an exercise to the reader.

5 Fuzzy closure systems

A natural generalisation of closure systems is to consider, instead of the set of closed elements, the degree to which each element is closed. For this aim, the equality between the element and its closure is replaced by the so-called \otimes -symmetric kernel relation, which is defined as $(a \approx_{\otimes} b) = \rho(a, b) \otimes \rho(b, a)$ for all $a, b \in A$. Thus, the degree in which an element $a \in A$ is closed is $(c(a) \approx_{\otimes} a)$ and, since every closure operator is inflationary, it is equal to $\rho(c(a), a)$. In [19], where the underlying structure is a Heyting algebra, the one-to-one relation between closure operators and

fuzzy closure systems is easily obtained following the ideas used for closure systems. However, this is not the case when the underlying structure is a general residuated lattice case.

Theorem 17. Let (A, ρ) be a complete fuzzy lattice and c be a closure operator on \mathbb{A} . For $\Phi_c \in L^A$, defined as $\Phi_c(a) = \rho(c(a), a)$, we have that $\text{Core}(\Phi_c) = \mathcal{F}_c$ is the closure system associated to c and, for all $x \in A$,

$$\prod (x^\rho \otimes \Phi_c) = \prod (x^\rho \cap \mathcal{F}_c) \quad (4)$$

In addition, for all $a \in A$, one has

$$\Phi_c(a) = \bigvee_{x \in \mathcal{F}_c} (a \approx_\otimes x) \quad (5)$$

Proof. Due to Theorem 7, to prove (4), it is enough to see that, for all $x, a \in A$,

$$\begin{aligned} & \bigwedge_{y \in A} ((\rho(x, y) \otimes \Phi_c(y)) \rightarrow \rho(a, y)) \\ &= \bigwedge_{s \in \mathcal{F}_c} (\rho(x, s) \rightarrow \rho(a, s)) \end{aligned}$$

On the one hand, the inequality \leq is straightforward because $\rho(x, s) \otimes \Phi_c(s) = \rho(x, s)$ for all $s \in \mathcal{F}_c$. On the other hand, for all $y \in A$,

$$\begin{aligned} & \rho(x, y) \otimes \Phi_c(y) \otimes (\rho(x, c(y)) \rightarrow \rho(a, c(y))) = \\ &= \rho(x, y) \otimes \rho(c(y), y) \otimes (\rho(x, c(y)) \rightarrow \rho(a, c(y))) \\ &= \rho(x, y) \otimes \rho(y, c(y)) \otimes \rho(c(y), y) \\ & \quad \otimes (\rho(x, c(y)) \rightarrow \rho(a, c(y))) \\ &\leq \rho(x, c(y)) \otimes \rho(c(y), y) \otimes (\rho(x, c(y)) \rightarrow \rho(a, c(y))) \\ &\leq \rho(a, c(y)) \otimes \rho(c(y), y) \leq \rho(a, y) \end{aligned}$$

Therefore, for all $a, x, y \in A$, we have that $c(y) \in \mathcal{F}_c$ and

$$\rho(x, c(y)) \rightarrow \rho(a, c(y)) \leq (\rho(x, y) \otimes \Phi_c(y)) \rightarrow \rho(a, y)$$

Finally, we prove (5). For all $a \in A$, For all $a \in A$ and $x \in \mathcal{F}_c$ (i.e. $c(x) = x$), by isotony of c and transitivity, we have that $(a \approx_\otimes x) = \rho(a, x) \otimes \rho(x, a) \leq \rho(c(a), c(x)) \otimes \rho(x, a) \leq \rho(c(a), a) = \Phi_c(a)$.

On the other hand, we have that $c(a) \in \mathcal{F}_c$ and $\Phi_c(a) = \rho(c(a), a) = (a \approx_\otimes c(a))$. \square

This leads to the definition of fuzzy closure systems.

Definition 18. Let (A, ρ) be a complete fuzzy lattice. We say that a fuzzy set $\Phi \in L^A$ is a fuzzy closure system if $\text{Core}(\Phi)$ is a closure system and, for all $a \in A$, the following condition holds

$$\Phi(a) = \bigvee_{x \in \text{Core}(\Phi)} (x \approx_\otimes a).$$

This additional property ensures that Φ is extensional.

Corollary 19. Let $\mathbb{A} = (A, \rho)$ be a complete fuzzy lattice. The following assertions hold:

- i. If Φ is a fuzzy closure system on \mathbb{A} , the mapping $c_\Phi: A \rightarrow A$ defined as $c_\Phi(a) = \prod (a^\rho \otimes \Phi)$ is a closure operator on \mathbb{A} .
- ii. If c is a closure operator on \mathbb{A} , the fuzzy set Φ_c defined as $\Phi_c(a) = \rho(c(a), a)$ is a fuzzy closure system on \mathbb{A} .
- iii. If Φ is a fuzzy closure system on \mathbb{A} , then $\Phi = \Phi_{c_\Phi}$.
- iv. If $c: A \rightarrow A$ is a closure operator on \mathbb{A} , then $c_{\Phi_c} = c$.

6 Conclusions and future work

In this paper, a version of the notion of closure system has been provided on general complete fuzzy lattices. This definition has been proved to behave properly in relation with closure operators, that is, there exists a bijection between both notions.

Moreover, the definition presented here has been compared with the original given by Bělohlávek, defined on L -powerset lattices, and have been proved to be equivalent. Hence, the former definition can be used in general complete fuzzy lattices avoiding the explicit use of L -powerset tools, namely the external operations given in (2).

As prospect of future work, a next step might be the extension of the concept of closure operator and consider it as some kind of fuzzy relation instead of a crisp mapping, in line with the work in [7], and study the relationship between fuzzy closure systems and closure relations.

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