

Correspondence Between Constructive Real Numbers and L Number in Constructive Mathematics

Xin Shen ^{1*}, Yiding Huang², Shixiang Wang³

¹Nanjing Foreign Language School, Nanjing, China
²Beijing 101 middle school, Beijing, China
³Oxford International College, Oxford, OX4 1BD, United Kingdom
 *Corresponding author. Email: 2830627975@qq.com

ABSTRACT

In a broad sense, construction real number is a constructable real number, relative to the set of some construction methods. This essay mainly focuses on the left numbers in constructive analysis. After introducing the basic context and conditions in Constructive Mathematics, connections between concepts are put into consideration and discovered with reference. Specific propositions are also discussed and utilized to simplify and prove the statements listed below. It is deduced that in the context of constructive mathematics, there exists an algorithm such that given any constructive real number, it produces the corresponding Left (L) number, given the presumption that with equivalent constructive real numbers (CRN), equivalent Left numbers can be the output. In addition, the Specker sequence is utilized to prove that the converse statement is false. It is hoped that the paper can introduce more people to basic concepts and application of theorems in analysis.

Keywords: Specker sequence, constructive mathematics, Specker Sequence

1. INTRODUCTION

Left number is a sequence of increasing rational numbers, which can be derived by some algorithm and is bounded above by a rational number. Every rational number can be thought of as a left number given as a geometric series of rational numbers. We focus on the correspondence between constructive real numbers and Left(L) numbers.

In this essay, some definitions in the constructive mathematical context are introduced and clarified including algorithm, constructive real number, decidable, and enumerable sets. Then, it is proved that there is an algorithm that given any constructive real number produces the corresponding left number. It is also

observed that the Specker sequence example shows that the converse statement is false. With further observation and deduction, it is logical to divide the proof of these two statements into three parts: 1) finding an algorithm using standard regulator to get an increasing sequence of rational numbers 2) proving that equivalent CRN give equivalent left numbers. 3) providing and explaining an example of a left number which does not converge to CRN. It is hoped that through research and proof of the above problems, the paper can provide some ideas for those who are interested in topology. In case that constructive mathematics is still unknown to most people, here are some information in the context of constructive mathematics to further enhance comprehension and understanding:

$(x \ll (U, \mathcal{M}))$	$\vdash \neg((x \ll (U, \mathcal{M})) \wedge (x \not\ll (U, \mathcal{M})))$
$x \in U$	$\vdash x \in U$
$x \in \mathcal{M}$	$\vdash (x \not\ll (U, \mathcal{M}))$
$(x \ll (U, \mathcal{M})) \wedge (U \subseteq V) \wedge (\mathcal{M} \subseteq \mathcal{N})$	$\vdash (x \ll (V, \mathcal{N}))$
$(x \not\ll (V, \mathcal{N})) \wedge (U \subseteq V) \wedge (\mathcal{M} \subseteq \mathcal{N})$	$\vdash (x \not\ll (U, \mathcal{M}))$
$(x \ll (U, \mathcal{M})) \wedge (x \not\ll (V, \mathcal{N}))$	$\vdash \exists y.(y \in U \cap \mathcal{N})$
	$\vdash (x \ll (A, \emptyset))$
	$\vdash \neg(x \not\ll (A, \emptyset))$
$(x \ll (U, \mathcal{M})) \wedge (x \ll (V, \mathcal{N}))$	$\vdash (x \ll (U \cap V, \mathcal{M} \cup \mathcal{N}))$
$(x \ll (U, \mathcal{M})) \wedge (x \not\ll (U \cap V, \mathcal{M} \cup \mathcal{N}))$	$\vdash (x \not\ll (V, \mathcal{N}))$
$(x \ll (U, \mathcal{M}))$	$\vdash (x \ll (\{y \mid y \ll (U, \mathcal{M})\}, \{y \mid y \not\ll (U, \mathcal{M})\}))$
$(x \not\ll (\{y \mid y \ll (U, \mathcal{M})\}, \{y \mid y \not\ll (U, \mathcal{M})\}))$	$\vdash (x \not\ll (U, \mathcal{M}))$

Figure 1 Linear logic for constructive mathematics

What is the relationship between constructive mathematics and classical mathematics? The relationship is clear as follows:

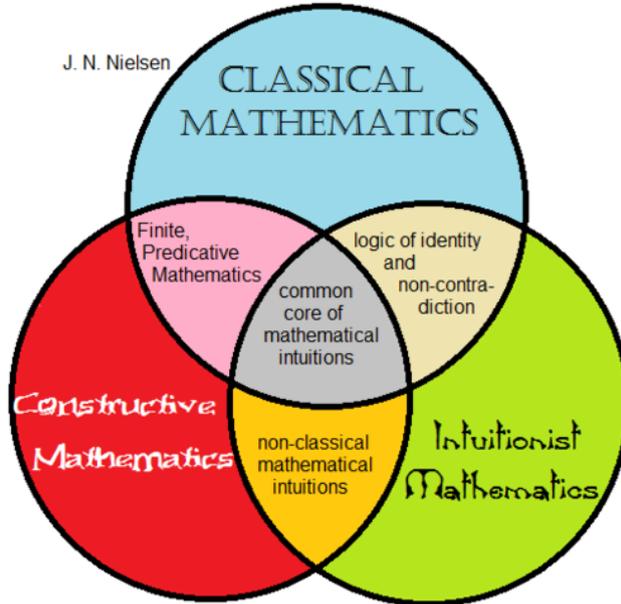


Figure 2 Is mathematics this whole structure – classicism, constructivism, and intuitionism taken together – plus whatever else is done in the name of mathematics?

2. DEFINITIONS

2.1. L(left) number

Left number is any strictly increasing algorithmically given sequence of rational numbers bounded from above by a bounding element θ , which is a rational number [1].

2.2. Computable (constructive) real number with Standard Regulator

There are two similar definitions that are equivalent:

There exists a computable function which, given any positive rational error bound ϵ , produces a rational number r such that $|r - a| \leq \epsilon$ [2].

There is a computable sequence of rational numbers q_i converging to a such that

$$|q_i - q_{i+1}| < 2^{-i}, \text{ for each } i \text{ [3].}$$

2.3 Computably Enumerable set

A set S of natural numbers is called computably enumerable if there is a partial computable function whose domain is exactly S , meaning that the function is defined if and only if its input is a member of S [4].

2.4 Decidability

There exists a general recursive function f such that $M = \{n: f(n) = 0\}$. In this case f is an algorithm for checking whether a natural number belongs to M :

indeed, $n \in M$ is equivalent to $f(n) = 0$. A decidable set of natural numbers is also often called a general recursive set or a recursive set.

2.5. Algorithm

In mathematics, an algorithm is a process that describes a series of steps that may be used to perform a mathematical computation.

3. PROOF

FIRST LET'S ASSUME THAT CONSTRUCTIVE REAL NUMBERS HAVE A STANDARD REGULATOR.

Let the CRN have a sequence $\alpha(n)$ equipped to the standard regulator. There exist an algorithm that transform the sequence into x , such that

$$x(n) = \alpha(n + 1) - \frac{1}{2^n} \quad (1)$$

In order for this sequence to be increasing, we should have:

$$\alpha(n + 1) - 2^{-n} < \alpha(n + 2) - 2^{-(n-1)} \quad (2)$$

which means

$$\alpha(n + 1) - \alpha(n + 2) < 2^{-(n+1)} \quad (3)$$

Now we prove that this is true.

From the definition of standard regulator, we know $\forall n, l \geq n, |\alpha(k) - \alpha(l)| < 2$ (4)
 Substitute n with n+1, k with n+1, and l with n+2, we have

$$|\alpha(n + 1) - \alpha(n + 2)| < 2^{-(n+1)} \quad (5)$$

which tells us that

$$\alpha(n + 1) - \alpha(n + 2) < 2^{-(n+1)} \quad (6)$$

Therefore, the sequence is strictly increasing.

The sequence is also bounded above by $\alpha(1) + \frac{1}{2}$. Therefore, we know that x is a left number.

Moreover, from the definition of CRN, which is provided above,

$$x(n + 1) - x(n) = \alpha(n + 1) - \alpha(n) + \frac{1}{2^{n+1}} \quad (7)$$

which is greater than 0.

So, from this we can generate a Left number x.

Let's first define the comparison between Left numbers:

if $L_1 < L_2$, then there exists a number in the sequence L_2 that is greater than every number in L_1

Moreover, let's define the equivalence of two L numbers:

if $L_1 \leq L_2$ and $L_1 \geq L_2$, then $L_1 = L_2$

To simplify, we use $\alpha(n)$ and $\alpha'(n)$ as two different sequences which converge are equivalent.

Let L number $x(n)$ be $\alpha(n) - 2^{-n}$ and L number $y(n)$ be $\alpha'(n) - 2^{-n}$

So we have

$$|\alpha(n) - \alpha'(n)| < 2^{-n} \quad (8)$$

$$s \quad |x(n) - y(n)| < 2^{-n} \quad (9)$$

Assume $x \leq y$ is false, then there exists a n such that for any i, $x(n) > y(i)$. Moreover, there exists a k, such that $k > n$ and $x(n + 1) - x(n) > 2^{-k}$.

Moreover, since x(n) and y(n) strictly increases, we get

$$|x(k) - y(k)| = x(k) - y(k) \geq x(n + 1) - y(k) > x(n + 1) - x(n) > 2^{-k} \quad (10)$$

Moreover, since

$$|x(n) - y(n)| < 2^{-n} \quad (11)$$

Substitute n with k, we have

$$|x(k) - y(k)| < 2^{-k} \quad (12)$$

This is a contradiction. Therefore, equivalent L numbers can be outputted when the input CRNs are equivalent.

To prove this, we can simply look for an example of left number that does not converge to CRN.

Let's first show that the problem is equivalent to finding as Specker sequence:

From the definition of Specker sequence, we know that the sequence is strictly increasing, has limit, and its supremum cannot be computed, which means that it's not a CRN.

Therefore, Specker sequence is Left number, and it has CRN which requires both a series and regulator.

Here is an example of the generating of Specker sequence [5]:

Consider recursively enumerable set of natural numbers, A, and it is not decidable, and a_i represents a kind of computable enumeration of A. It does not contain any form of repetition. Next, let's define a sequence of rational numbers, call it q_n .

$$q_n = \sum_{i=0}^n 2^{-a_i-1} \quad (13)$$

Each q_n is rational and nonnegative, and as a sequence, q_n is strictly increasing. Moreover, estimating each q_n against the series is reasonable, because the kind of computable enumeration of A has no repetition.

$$\sum_{i=0}^{\infty} 2^{-i-1} = 1 \quad (14)$$

Therefore, it is clear that the sequence q_n is bounded above by 1, a rational number.

With this in mind, we know that q_n has a supreme, satisfying one of the key propositions of Specker Sequence.

Plus, since it is known that x is not a computable real number, if x were computable then let's assume that there existed a computable function r(n) such that $|q_j - q_i| < \frac{1}{n}$, with condition $i, j > r(n)$. After comparing the binary expansion of x with the binary expansion of q_i for increasing values of I, it is logical to compute r. Consider the definition of q_i . Each time i increases by 1, there is a single binary digit that goes from 0 to 1. See the pattern below.

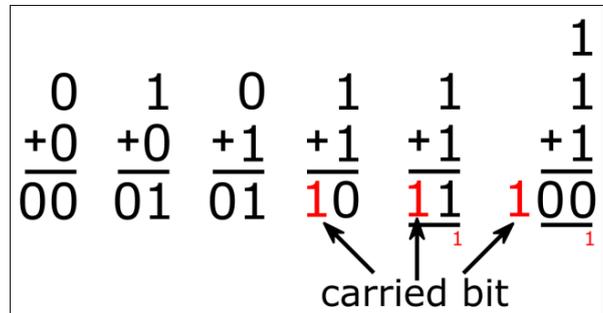


Figure 3 Property of binary numbers

There will be, therefore, an n such that a very large initial segment of x is already judged by q_n . No additional binary digits in that part could be transformed

