# Improvement of Gradient Projection Algorithm for Nonlinear Programming 

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$$
M A=\left(\begin{array}{cccc}
* & * & \cdots & *  \tag{2}\\
0 & * & \cdots & * \\
0 & 0 & \ddots & * \\
0 & 0 & \cdots & * \\
0 & 0 & \cdots & 0 \\
0 & 0 & \cdots & 0 \\
0 & 0 & \cdots & 0
\end{array}\right)
$$

The (2) both sides of transposition, obtain
$A^{T} M^{T}$

$$
=\left(\begin{array}{ccccccc}
* & 0 & 0 & 0 & 0 & \cdots & 0  \tag{3}\\
* & * & 0 & 0 & 0 & \cdots & 0 \\
\vdots & \vdots & \ddots & \vdots & 0 & \cdots & 0 \\
* & * & * & * & 0 & \cdots & 0
\end{array}\right)
$$

Inside $M_{1}, \mathrm{M}_{2}, \cdots, M_{n}$ are lower triangular matrix, so

$$
M^{T}=\left(M_{n} \cdots M_{2} M_{1}\right)^{T}=M_{1}{ }^{T} M_{2}{ }^{T} \cdots M_{n}{ }^{T}
$$

are upper triangular matrix.
Theorem 1.
Assume $\bar{U}=\left(u_{n+1}, u_{n+2}, \cdots, u_{m}\right)$, then $\bar{U}$ is a group of base on vertical subspace of A.

Proof.
Assume

$$
\begin{gathered}
M^{T}=\left(u_{1}, u_{2}, \cdots, u_{n}, \cdots, u_{m}\right), \\
\text { so (3) can be changed to } \\
A^{T}\left(u_{1}, u_{2}, \cdots, u_{n}, \cdots, u_{m}\right) \\
=\left(\begin{array}{cccccccc}
* & 0 & 0 & 0 & 0 & \cdots & 0 \\
* & * & 0 & 0 & 0 & \cdots & 0 \\
\vdots & \vdots & \ddots & \vdots & 0 & \cdots & 0 \\
* & * & * & * & 0 & \cdots & 0
\end{array}\right)
\end{gathered}
$$

Because the rank of $A$ is $n$,so the matrix

$$
\left(\begin{array}{ccccccc}
* & 0 & 0 & 0 & 0 & \cdots & 0 \\
* & * & 0 & 0 & 0 & \cdots & 0 \\
\vdots & \vdots & \ddots & \vdots & 0 & \cdots & 0 \\
* & * & * & * & 0 & \cdots & 0
\end{array}\right)
$$

From the $(\mathrm{n}+1)$ th column is vector 0 . so

$$
\begin{aligned}
& A^{T} u_{n+1}=0 \\
& A^{T} u_{m}=0
\end{aligned}
$$

$\bar{U}$ is a group of base on vertical subspace of A.
According to Theorem 1, we can obtain the corresponding algorithm to solve $\bar{U}$, denoted Algorithm 1.

Because the effective constraint coefficient matrix of large-scale problems often have some sparse structures, we can make the constraint matrix become a banded structure by adjusting the operative constraint and the order of unknown sequence, so using algorithm 1 calculation $\bar{U}$, it can not only looser the sparsity of the constraint matrix requirements, but it can also decrease the storage capacity of new method than the usual LU decomposition.

## III. The gradient projection algorithm for NONLINEAR PROGRAMMING

Considering the nonlinear programming problem $\min \{f(x) \mid x \in S\}$,
$S$ is a set defined by

$$
\begin{aligned}
& S=\{x \mid A x=b, x \geq 0\}, \\
& S_{+}=\{x \mid x \in S, x \succ 0\}
\end{aligned}
$$

Algorithm 2
Step 1. Set the initial point $x^{0} \in S_{+}, k=0$.
Step 2. Using $x^{k}$ the definition of diagonal matrix $X=\operatorname{diag}\left(x_{1}^{k}, \cdots, x_{n}^{k}\right)$,
do transform

$$
\begin{gathered}
\bar{x}=X^{-1} x, \\
\bar{c}=X C, \\
\bar{A}=A X
\end{gathered}
$$

Step 3. Using algorithm 1 to solve $u$, $u$ satisfy: $\bar{A}^{T} u=0$.

Step 4. Solve the gradient projection direction:

$$
\begin{gathered}
d f=d f G * \bar{x}+d f b \\
\bar{c}_{p}=-u u^{T} d f \\
\bar{d}=\frac{1}{\left\|\bar{c}_{p}\right\|} \bar{c}_{p}
\end{gathered}
$$

Step 5. Iterative:

$$
\bar{x}^{-k+1}=e+\alpha \bar{d}, 0<\alpha<1,
$$

$$
x^{k+1}=X^{-k+1}
$$

Step 6. If the termination condition is satisfied, then the end; otherwise $k:=k+1$, return step 2 .

## IV. EXAMPLE

Example 1:
solve the nonlinear programming

$$
\begin{aligned}
& \min \left(x_{2}^{2}+x_{4}^{2}+\cdots+x_{98}^{2}+x_{100}^{2}\right) \\
& \quad \text { s.t. } x_{i}+x_{i+1}=0.1, i=1,2, \cdots, 100 \\
& x_{i} \geq 0, i=1,2, \cdots, 101
\end{aligned}
$$

Obviously, the optimal objective function value of the linear programming is 0 , one of it is:

$$
\begin{aligned}
x_{2 i+1} & =0.1, i=0,1,2, \cdots, 50 \\
x_{2 i} & =0.0, i=1,2, \cdots, 50 .
\end{aligned}
$$

From the initial point of

$$
x_{1}=x_{2}=\cdots=x_{100}=x_{101}=0.05
$$

after 311 iterations, the optimalvalue of 4.8640e-013.
Example 2:
solve the nonlinear programming

$$
\begin{gathered}
\min \left(x_{3}^{2}+x_{6}^{2}+\cdots+x_{99}^{2}+x_{102}^{2}\right) \\
\text { s.t. } x_{i}+x_{i+1}+x_{i+2}=0.6 \\
i=1,2, \cdots, 100 \\
\quad x_{i} \geq 0, i=1,2, \cdots, 102
\end{gathered}
$$

Obviously, the optimal objective function value of the linear programming is 0 , one of it is:

$$
x_{3 i+1}=0.4 \quad, \quad x_{3 i+2}=0.2 \quad, \quad x_{3 i+3}=0.0
$$ $i=1,2, \cdots, 33$.

From the initial point of

$$
x_{1}=x_{2}=\cdots=x_{100}=x_{101}=0.2
$$

after 306 iterations, the optimalvalue of 4.7609e-013.

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