

The structure of solution sets of fuzzy relation equations

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Abstract

In this paper, we consider the structure of solution sets of fuzzy relation equations over complete Boolean algebras. We show that each solution of a system of fuzzy relation equations can be represented by a linear combination of a special solution of its and some certain solutions of the homogeneous equations associated with the system.

Keywords: Fuzzy relation equations; Complete Boolean algebra; Structure of solution set

1. Introduction

System of fuzzy relation equations is an important topic in fuzzy set theory. In 1976, Sanchez [21] first introduced fuzzy relation equations with sup-inf composition in complete Brouwerian lattices. Since then, many authors investigated the methods for solving fuzzy relation equations with different composite operators over various special Brouwerian lattices. Among them, for finite fuzzy relation equations with sup-inf composition, Higashi et al. [11] showed that the solution set can be determined by minimal solutions and the greatest solution in the linear lattice $[0,1]$. Zhao [27] determined its entirely solution set in complete and completely distributive lattices. De Baets [4] determined all minimal solutions and Wang [24] gave a formula of the number of minimal solutions under the condition that all elements of its right-hand side have irredundant finite join-irreducible decompositions. And many other works about solving fuzzy relation equations were published (see [1, 5, 8, 10, 12, 13, 14, 15, 16, 17, 19, 20, 22, 23]). Compared with the linear algebraic systems, the main difference with fuzzy relation equations is the operations and the fundamental domain. Recently, there are some ideas which learn from linear algebraic systems for studying fuzzy relation equations. In 2004, Perfilieva [18] introduced the concept of semi-linear space based on BL-algebra. Di Nola et al. [7] introduced the concepts of semi-linear space, basis and linear independent based on MV-algebra, and got to connect the semi-linear space and systems of fuzzy relation equations. Wang and Zhao[26] described the solution set of systems

of fuzzy relation equations through n -dimensional vectors over bounded Brouwerian lattices.

A generic version of fuzzy relation equation arises in the form

$$A \odot \mathbf{x} = \mathbf{b} \quad (1)$$

or

$$\bigvee_{j \in J} (a_{ij} \wedge x_j) = b_i, \text{ for all } i \in I,$$

where $A = (a_{ij})_{(i,j) \in I \times J}$ and $\mathbf{b} = (b_i)_{i \in I}^T$ are known (I and J are index sets with $I = \{1, 2, \dots, m\}$ and $J = \{1, 2, \dots, n\}$), $\mathbf{x} = (x_j)_{j \in J}^T$ is unknown, \odot denotes the sup-inf composite operation, and all b_i, x_j, a_{ij} 's are in a complete Brouwerian lattice L . An \mathbf{x} which satisfies Eq.(1) is called a solution of (1), the solution set of Eq.(1) is denoted by $\mathcal{X} = \{\mathbf{x} : A \odot \mathbf{x} = \mathbf{b}\}$. If $\mathcal{X} \neq \emptyset$, the greatest element of \mathcal{X} is denoted by $\mathbf{x}^* = (x_j^*)_{j \in J}^T$, and we denote $\mathcal{X}^0 = \{\mathbf{x} \in \mathcal{X} : \mathbf{x} \text{ is a minimal element of } \mathcal{X}\}$. The following equation is called the homogeneous equations associated with Eq.(1)

$$A \odot \mathbf{x} = \mathbf{0}, \quad (2)$$

where $\mathbf{0} = (0, 0, \dots, 0)^T$ is a zero vector. \mathcal{Y} and \mathbf{y}^* denote the solution set and the greatest solution of Eq.(2), respectively.

In this paper, we consider the structure of Eq.(1) over a special complete Brouwerian lattice, that is, a complete Boolean algebra $(L, \vee, \wedge, 0, 1)$. We will describe the solution set of Eq.(1) by a linear combination of a special solution of its and some certain solutions of Eq.(2). This paper is organized as follows. Section 2 presents some preliminary information which includes fuzzy relation equations and lattices. Section 3 gives some properties of elements in complete Boolean algebras. In section 4, we describe the structure of solution sets of fuzzy relation equations. Conclusions are given in section 5.

2. Preliminaries

In this section, we give some definitions and previous results for the sake of convenience.

Definition 2.1 [2] *A lattice L is Brouwerian if for any pair of elements $a, b \in L$, the greatest element*

$x \in L$, denoted by aob , satisfying the inequality $a \wedge x \leq b$ exists.

Definition 2.2 [3, 23] Let L be a complete lattice and $a \in L$. If T is a finite subset of L and S (finite or infinite) is a subset of L , the representation $a = \bigvee T$ is called an irredundant finite join decomposition of a if $a \neq \bigvee_{t \in T} (T - \{t\})$ for each $t \in T$, and the representation $a = \bigvee S$ is called a minimal join decomposition of a if $a \neq q \vee \bigvee_{s \in S, s \neq p} s$ for each $p \in S$ and for any $q < p$.

Let L be a complete lattice. For each $a \in L$, denote $\mathcal{M}_a = \{M \subseteq L : a = \bigvee M \text{ is a minimal join decomposition of } a\}$. It is clear that $\{a\} \in \mathcal{M}_a$.

Definition 2.3 [3] Let L be a lattice with a least element 0 and a greatest element 1. A complement of an element $a \in L$ is an element $b \in L$ such that $a \wedge b = 0$ and $a \vee b = 1$. We write $b = a'$. L is called complemented if each of its elements has a complement, and L is relatively complemented if the quotient sublattice $a/b = \{x \in L : b \leq x \leq a\}$ is complemented for each pair of elements $a > b$ in L . A complemented, distributive lattice is called a Boolean algebra. A complete Boolean algebra L is infinitely distributive; i.e., if $a \in L$ and S is a subset of L , then $a \wedge \bigvee S = \bigvee_{s \in S} (a \wedge s)$.

Let L be a complete relatively complemented lattice, $a \in L$ and $a \neq 0$, $b \leq a$, we denote the complement of b in the sublattice $a/0$ by b'_a .

Lemma 2.1 [3] Every complemented modular lattice is relatively complemented. In particular, Boolean algebras are relatively complemented.

Lemma 2.2 [2] Let L be a Boolean algebra. Then $(a \wedge b)' = a' \vee b'$ for all $a, b \in L$.

Lemma 2.3 [9] Let L be a Boolean algebra. Then $aob = a' \vee b$ for all $a, b \in L$.

Lemma 2.4 [21] If $\mathcal{X} \neq \emptyset$, then $\mathbf{x}^* = (\bigwedge_{i \in I} (a_{ij} ob_i))_{j \in J}^T$ is the greatest element of \mathcal{X} . If $\mathbf{x}_1, \mathbf{x}_2 \in \mathcal{X}$ and $\mathbf{x}_1 \leq \mathbf{x}_2$, then $\mathbf{x} \in \mathcal{X}$ for any \mathbf{x} such that $\mathbf{x}_1 \leq \mathbf{x} \leq \mathbf{x}_2$.

Lemma 2.5 [27] For equation $\bigvee_{j \in J} (a_j \wedge x_j) = b$, $\mathcal{X} \neq \emptyset$ if and only if $b \leq \bigvee_{j \in J} a_j$.

Remark 2.1 For equation $\bigvee_{j \in J} (a_j \wedge x_j) = b$, by Lemma 2.5, $\mathcal{X} \neq \emptyset$ if and only if $\mathbf{x} = (b)_{j \in J} \in \mathcal{X}$.

Lemma 2.6 [25] For equation $\bigvee_{j \in J} (a_j \wedge x_j) = b$, if $\mathbf{x}_* = (x_{j*})_{j \in J}^T$ is a minimal element of \mathcal{X} , then $b = \bigvee_{j \in J} x_{j*}$, and for all $j \in J$, $x_{j*} \leq a_j$.

3. Properties of elements in complete Boolean algebras

In this section, we focus on discussing the properties of elements in complete Boolean algebras. In the rest of this paper, unless otherwise specified, we always assume that L is a complete Boolean algebra with the least element 0 and the greatest element 1. For each $a \in L$, $a \neq 0$, denote $D(a) = \{x \leq a : \text{there is some element } y < a \text{ such that } x \vee y = a\}$.

Theorem 3.1 Let $a, b \in L$, and $a \leq b$. Then $a'_b = a' \wedge b$.

Proof. Since $a \vee a' = 1$ and $a \wedge a' = 0$, it follows that $a \vee (a' \wedge b) = (a \vee a') \wedge (a \vee b) = 1 \wedge b = b$, and $a \wedge (a' \wedge b) = (a \wedge a') \wedge b = 0 \wedge b = 0$. Then $a'_b = a' \wedge b$. ■

Theorem 3.2 Let $a \in L$. Then for every $b \in D(a)$, there is some $M \in \mathcal{M}_a$ such that $b \in M$.

Proof. If $b = a$, then $b \in \{a\}$ and $\{a\} \in \mathcal{M}_a$. If $b < a$, then there is b'_a such that $b \vee b'_a = a$ and $b \wedge b'_a = 0$. We infer that $M = \{b, b'_a\} \in \mathcal{M}_a$. If there exists $c < b$ such that $c \vee b'_a = a$, then $c \wedge b'_a = 0$ since $b \wedge b'_a = 0$. This implies that c is a relatively complement of b'_a . So b'_a has two different complements, a contradiction. Similarly, there is no element $c < b'_a$ such that $b \vee c = a$. Therefore, by Definition 2.2, $M = \{b, b'_a\} \in \mathcal{M}_a$. ■

In the following two theorems, the index set J is defined as previous, that is, $J = \{1, 2, \dots, n\}$.

Theorem 3.3 If $a \in L$ and $a = \bigvee_{j \in J} t_j$ is an irredundant finite join decomposition of a , then there exists $M = \{s_j \in L : j \in J\} \in \mathcal{M}_a$ such that $s_j \leq t_j$ for all $j \in J$.

Proof. Choose any $k \in J$, let $a_k = \bigvee_{j \in J, j \neq k} t_j$, and $S = \{x \in L : x \vee a_k = a\}$. Certainly, $S \neq \emptyset$ since $t_k \in S$. Because of the upper continuity of L , there is a minimal element $s_k \in S$ (With the distributivity, in fact, s_k is the least element of S), that is, $a = s_k \vee a_k$ and $a \neq x \vee a_k$ for any $x < s_k$. Next, choose any $k_1 \in J - \{k\}$ and let $a_{k_1} = s_k \vee \bigvee_{j \in J, j \neq k, k_1} t_j$. It is similar that there is an element $s_{k_1} \in L$ such that $a = s_{k_1} \vee s_k \vee a_{k_1}$ and $a \neq x \vee s_k \vee a_{k_1}$ for any $x < s_{k_1}$. Continuing in this way we can obtain a set $M = \{s_j \in L : j \in J\} \in \mathcal{M}_a$ such that $s_j \leq t_j$ for all $j \in J$. ■

Theorem 3.4 Let $a, b \in L$, $b < a$ and $b \in D(a)$. If $a = b \vee \bigvee_{j \in J} t_j$, then there exist elements $s_j \leq t_j$, $j \in J$ such that $a = b \vee \bigvee_{j \in J} s_j$ and $a \neq \bigvee_{j \in J} s_j$. Moreover, there is an element $M \in \mathcal{M}_a$ such that $b \in M$ and for every $u \in M$, $u \neq b$, there exists some t_j such that $u \leq t_j$.

Proof. Assume that $b < a$ and $b \in D(a)$, b'_a is the relatively complement of b in sublattice $a/0$. Then $b \vee b'_a = a$. Let $s_j = t_j \wedge b'_a$, $j \in J$. Since $a = b \vee b'_a$

and $a = b \vee \bigvee_{j \in J} t_j$, we have $a = b \vee (b'_a \wedge \bigvee_{j \in J} t_j) = b \vee \bigvee_{j \in J} (b'_a \wedge t_j) = b \vee \bigvee_{j \in J} s_j$. In addition, with the fact $b'_a < a$, we have $\bigvee_{j \in J} s_j = \bigvee_{j \in J} (b'_a \wedge t_j) \leq b'_a < a$. Hence $a \neq \bigvee_{j \in J} s_j$. Now removing the redundant elements from the representation $a = b \vee \bigvee_{j \in J} s_j$, we can get an irredundant representation $a = b \vee \bigvee_{j \in J'} s_j$ and $a \neq b \vee \bigvee_{j \in J', j \neq k} s_j$ for each $k \in J'$, where $J' \subseteq J$. Similar to the proof of Theorem 3.2, there is no element $c < b$ such that $a = c \vee \bigvee_{j \in J'} s_j$. Furthermore, according to Theorem 3.3, there are elements $u_j \in L$, $j \in J'$, and $u_j \leq s_j$ for all $j \in J'$ such that $M = \{u_j \in L : j \in J'\} \cup \{b\} \in \mathcal{M}_a$. It is clear that $u_j \leq s_j \leq t_j$ for all $j \in J'$, the proof is complete. ■

4. Structure of solution set of fuzzy relation equations

In this section, we will discuss the structure of solution set of Eq.(1) similar to linear algebraic systems. In the following, for the sake of convenience, we sometimes denote $a \vee b$ and $a \wedge b$ by $a + b$ and ab , respectively.

For Eq.(2), it is clear that $\mathcal{Y} \neq \emptyset$ since $\mathbf{0}_{n \times 1} = (0, 0, \dots, 0)^T \in \mathcal{Y}$, the greatest solution is $\mathbf{y}^* = (y_j^*)_{j \in J}^T = (\bigwedge_{i \in I} a_{ij})_{j \in J}^T = ((\bigvee_{i \in I} a_{ij})')_{j \in J}^T$, and $\mathbf{y}^* \leq \mathbf{x}^*$. Now for each $k \in J$, define $\mathbf{e}_k = (y_j)_{j \in J}^T$ with $y_k = (\bigvee_{i \in I} a_{ij})'$ and $y_j = 0$ if $j \neq k$. Then $\mathbf{e}_k \in \mathcal{Y}$ for all $k \in J$, and for any $\mathbf{y} = (y_j)_{j \in J}^T \in \mathcal{Y}$, we have $\mathbf{y} = y_1 \mathbf{e}_1 + y_2 \mathbf{e}_2 + \dots + y_n \mathbf{e}_n$. Therefore, $\mathcal{Y} = \{y_1 \mathbf{e}_1 + y_2 \mathbf{e}_2 + \dots + y_n \mathbf{e}_n : y_j \leq y_j^* \text{ for all } j \in J\}$.

Now we consider the structure of solution set of Eq.(1).

Theorem 4.1 *If $\mathbf{x} \in \mathcal{X}$ and $\mathbf{y} \in \mathcal{Y}$, then $\mathbf{x} \vee \mathbf{y} \in \mathcal{X}$ and $\mathbf{x} \wedge \mathbf{y} \in \mathcal{Y}$.*

Proof. Suppose that $\mathbf{x} = (x_j)_{j \in J}^T \in \mathcal{X}$ and $\mathbf{y} = (y_j)_{j \in J}^T \in \mathcal{Y}$. Then $\bigvee_{j \in J} (a_{ij} \wedge x_j) = b_i$ and $\bigvee_{j \in J} (a_{ij} \wedge y_j) = 0$ for all $i \in I$. Thus for all $i \in I$, $\bigvee_{j \in J} (a_{ij} \wedge (x_j \vee y_j)) = \bigvee_{j \in J} ((a_{ij} \wedge x_j) \vee (a_{ij} \wedge y_j)) = \bigvee_{j \in J} (a_{ij} \wedge x_j) \vee \bigvee_{j \in J} (a_{ij} \wedge y_j) = b_i \vee 0 = b_i$, and $\bigvee_{j \in J} (a_{ij} \wedge (x_j \wedge y_j)) \leq \bigvee_{j \in J} (a_{ij} \wedge y_j) = 0$. This implies that $\bigvee_{j \in J} (a_{ij} \wedge (x_j \wedge y_j)) = 0$. Therefore, $\mathbf{x} \vee \mathbf{y} \in \mathcal{X}$ and $\mathbf{x} \wedge \mathbf{y} \in \mathcal{Y}$. ■

Theorem 4.2 *For any $\mathbf{x} \in \mathcal{X}$, there is a solution $\mathbf{w} \in \mathcal{X}$ such that $\mathbf{x} = \mathbf{w} \vee (\mathbf{y}^* \wedge \mathbf{x})$.*

Proof. Suppose that $\mathbf{x} = (x_j)_{j \in J}^T \in \mathcal{X}$. Then $\mathbf{y}^* \wedge \mathbf{x} = ((\bigvee_{i \in I} a_{ij})' \wedge x_j)_{j \in J}^T$. Let $\mathbf{w} = (w_j)_{j \in J}^T$, where $w_j = ((\bigvee_{i \in I} a_{ij})' \wedge x_j)_{x_j}$ for all $j \in J$. Since $((\bigvee_{i \in I} a_{ij})' \wedge x_j)_{x_j} = ((\bigvee_{i \in I} a_{ij})' \wedge x_j)' \wedge x_j = ((\bigvee_{i \in I} a_{ij}) \vee x_j') \wedge x_j = ((\bigvee_{i \in I} a_{ij}) \wedge x_j) \vee (x_j \wedge x_j') = ((\bigvee_{i \in I} a_{ij}) \wedge x_j) \vee 0 = (\bigvee_{i \in I} a_{ij}) \wedge x_j$, we have $\bigvee_{j \in J} (a_{ij} \wedge w_j) = \bigvee_{j \in J} (a_{ij} \wedge ((\bigvee_{i \in I} a_{ij}) \wedge x_j)) = \bigvee_{j \in J} (a_{ij} \wedge x_j) = b_i$ for all $i \in I$. Thus $\mathbf{w} \in \mathcal{X}$. It can be easily shown that $\mathbf{x} = \mathbf{w} \vee (\mathbf{y}^* \wedge \mathbf{x})$. ■

Remark 4.1 *For any $\mathbf{x} = (x_j)_{j \in J}^T \in \mathcal{X}$, by Theorem 4.1, $\mathbf{y}^* \wedge \mathbf{x} = (y_j)_{j \in J}^T \in \mathcal{Y}$ holds. It follows that $\mathbf{y}^* \wedge \mathbf{x} = y_1 \mathbf{e}_1 + y_2 \mathbf{e}_2 + \dots + y_n \mathbf{e}_n$. Consequently, according to Theorem 4.2, there is a special solution $\mathbf{w} \in \mathcal{X}$ such that $\mathbf{x} = \mathbf{w} + y_1 \mathbf{e}_1 + y_2 \mathbf{e}_2 + \dots + y_n \mathbf{e}_n$. But if $\mathbf{x}_1, \mathbf{x}_2 \in \mathcal{X}$, and $\mathbf{x}_1 \neq \mathbf{x}_2$, then the special solution $\mathbf{w} \in \mathcal{X}$ need not be the same one.*

Let $\mathbf{y}^{*'} = (y_j^{*'})_{j \in J}^T = (\bigvee_{i \in I} a_{ij})_{j \in J}^T$, and $\mathbf{w}_0 = \mathbf{x}^* \wedge \mathbf{y}^{*'}$. Since for all $i \in I$, $\bigvee_{j \in J} (a_{ij} \wedge (x_j^* \wedge (\bigvee_{i \in I} a_{ij}))) = \bigvee_{j \in J} (a_{ij} \wedge x_j^*) = b_i$, this implies that $\mathbf{w}_0 \in \mathcal{X}$. Furthermore, we have the following conclusions.

Theorem 4.3 *For any $\mathbf{x} \in \mathcal{X}$, $\mathbf{x} \wedge \mathbf{w}_0 \in \mathcal{X}$.*

Proof. Suppose that $\mathbf{x} = (x_j)_{j \in J}^T \in \mathcal{X}$. Since for all $i \in I$, $\bigvee_{j \in J} (a_{ij} \wedge x_j \wedge (x_j^* \wedge \bigvee_{i \in I} a_{ij})) = \bigvee_{j \in J} (a_{ij} \wedge x_j) = b_i$, it follows that $\mathbf{x} \wedge \mathbf{w}_0 \in \mathcal{X}$. ■

Theorem 4.4 *For any $\mathbf{x} \in \mathcal{X}^0$, $\mathbf{x} \leq \mathbf{y}^{*'}$.*

Proof. According to Theorem 4.3, for any $\mathbf{x} \in \mathcal{X}^0$, we have $\mathbf{x} \wedge \mathbf{w}_0 \in \mathcal{X}$, then $\mathbf{x} \wedge \mathbf{w}_0 = \mathbf{x} \wedge \mathbf{x}^* \wedge \mathbf{y}^{*'}$ = $\mathbf{x} \wedge \mathbf{y}^{*'}$ $\in \mathcal{X}$. Since $\mathbf{x} \wedge \mathbf{y}^{*'}$ $\leq \mathbf{x}$ and $\mathbf{x} \in \mathcal{X}^0$, this implies that $\mathbf{x} \wedge \mathbf{y}^{*'}$ = \mathbf{x} . Consequently, $\mathbf{x} \leq \mathbf{y}^{*'}$. ■

Theorem 4.5 *If $\mathbf{x} \in \mathcal{X}$, $\mathbf{x} \wedge \mathbf{y}^* = \mathbf{0}$, then there are no elements $\mathbf{x}' \in \mathcal{X}$, $\mathbf{x}' \neq \mathbf{x}$, and $\mathbf{y} \in \mathcal{Y}$ such that $\mathbf{x} = \mathbf{x}' \vee \mathbf{y}$.*

Proof. If there exist elements $\mathbf{x}' \in \mathcal{X}$, $\mathbf{x}' \neq \mathbf{x}$, and $\mathbf{y} \in \mathcal{Y}$ such that $\mathbf{x} = \mathbf{x}' \vee \mathbf{y}$, then $\mathbf{y} < \mathbf{x}$, and $\mathbf{y} \leq \mathbf{x} \wedge \mathbf{y}^* = \mathbf{0}$. Hence $\mathbf{y} = \mathbf{0}$. Therefore, $\mathbf{x} = \mathbf{x}'$, a contradiction. ■

Theorem 4.6 *If $\mathbf{x} \in \mathcal{X}$, then $\mathbf{x} \wedge \mathbf{y}^* = \mathbf{0}$ if and only if $\mathbf{x} \leq \mathbf{w}_0$.*

Proof. If $\mathbf{x} \wedge \mathbf{y}^* = \mathbf{0}$, then $\mathbf{x} = \mathbf{x} \wedge \mathbf{1} = \mathbf{x} \wedge (\mathbf{y}^* \vee \mathbf{y}^{*'}) = (\mathbf{x} \wedge \mathbf{y}^*) \vee (\mathbf{x} \wedge \mathbf{y}^{*'}) = \mathbf{0} \vee (\mathbf{x} \wedge \mathbf{y}^{*'}) = \mathbf{x} \wedge \mathbf{y}^{*'}$. Thus $\mathbf{x} \leq \mathbf{y}^{*'}$. Hence $\mathbf{x} \leq \mathbf{x}^* \wedge \mathbf{y}^{*'}$ = \mathbf{w}_0 . Conversely, suppose that $\mathbf{x} \leq \mathbf{w}_0$, then $\mathbf{x} \wedge \mathbf{y}^* \leq \mathbf{w}_0 \wedge \mathbf{y}^* = (\mathbf{x}^* \wedge \mathbf{y}^{*'}) \wedge \mathbf{y}^* = \mathbf{y}^{*' \wedge} \mathbf{y}^* = \mathbf{0}$. Therefore, $\mathbf{x} \wedge \mathbf{y}^* = \mathbf{0}$. ■

Theorem 4.7 *Let $\mathbf{x} \in \mathcal{X}$. If $\mathbf{x} \geq \mathbf{w}_0$, then there is an element $\mathbf{y} \in \mathcal{Y}$ such that $\mathbf{x} = \mathbf{w}_0 \vee \mathbf{y}$.*

Proof. Suppose that $\mathbf{x} \in \mathcal{X}$ and $\mathbf{x} \geq \mathbf{w}_0$. Let $\mathbf{y} = \mathbf{x} \wedge \mathbf{y}^*$. Then $\mathbf{y} \in \mathcal{Y}$, and $\mathbf{w}_0 \vee \mathbf{y} = \mathbf{w}_0 \vee (\mathbf{x} \wedge \mathbf{y}^*) = (\mathbf{w}_0 \vee \mathbf{x}) \wedge (\mathbf{w}_0 \vee \mathbf{y}^*) = \mathbf{x} \wedge ((\mathbf{x}^* \wedge \mathbf{y}^{*'}) \vee \mathbf{y}^*) = \mathbf{x} \wedge ((\mathbf{x}^* \vee \mathbf{y}^*) \wedge (\mathbf{y}^{*' \vee} \mathbf{y}^*)) = \mathbf{x} \wedge (\mathbf{x}^* \wedge \mathbf{1}) = \mathbf{x}$. ■

Remark 4.2 *Let $\mathcal{X}_0 = \{\mathbf{x} \in \mathcal{X} : \mathbf{x} \geq \mathbf{w}_0\}$. By Theorems 4.1 and 4.7, we have $\mathcal{X}_0 = \{\mathbf{w}_0 + y_1 \mathbf{e}_1 + y_2 \mathbf{e}_2 + \dots + y_n \mathbf{e}_n : (y_1, y_2, \dots, y_n)^T \in \mathcal{Y}\}$. But according to Theorems 4.5 and 4.6, if $\mathbf{x} \in \mathcal{X}$, and $\mathbf{x} < \mathbf{w}_0$, then \mathbf{x} can be represented only by the combination of itself and zero element.*

Example 4.1 Let L be the lattice of power set of set $A = \{1, 2, 3, 4\}$. Denote $1 = A$, $0 = \emptyset$, $a_1 = \{1, 2, 3\}$, $a_2 = \{1, 2, 4\}$, $a_3 = \{1, 3, 4\}$, $a_4 = \{2, 3, 4\}$, $b_1 = \{1, 2\}$, $b_2 = \{1, 3\}$, $b_3 = \{1, 4\}$, $b_4 = \{2, 3\}$, $b_5 = \{2, 4\}$, $b_6 = \{3, 4\}$, $c_1 = \{1\}$, $c_2 = \{2\}$, $c_3 = \{3\}$, $c_4 = \{4\}$. Consider equation $A \odot \mathbf{x} = \mathbf{b}$, where

$$A = \begin{pmatrix} c_1 & b_3 & b_4 & b_1 \\ b_2 & b_1 & b_5 & b_6 \\ a_1 & b_5 & c_4 & b_3 \end{pmatrix},$$

$\mathbf{b} = (c_1, b_2, a_1)^T$. Obviously, $\mathcal{X} \neq \emptyset$, and $\mathbf{x}^* = (1, b_2, c_1, b_2)^T$, $\mathcal{X}^0 = \{\mathbf{x}_{*1} = (a_1, 0, 0, 0)^T, \mathbf{x}_{*2} = (b_4, c_1, 0, c_1)^T\}$. The associated equation is $A \odot \mathbf{x} = \mathbf{0}$, and $\mathbf{y}^* = (c_4, c_3, c_1, 0)^T$, $\mathbf{e}_1 = (c_4, 0, 0, 0)^T$, $\mathbf{e}_2 = (0, c_3, 0, 0)^T$, $\mathbf{e}_3 = (0, 0, c_1, 0)^T$, $\mathbf{e}_4 = (0, 0, 0, 0)^T$. Consider a solution $\mathbf{x}_1 = (a_1, c_1, c_1, 0)^T \in \mathcal{X}$, since $\mathbf{y}^* \wedge \mathbf{x} = (0, 0, c_1, 0)^T$, we have $\mathbf{x}_1 = \mathbf{w} + 0\mathbf{e}_1 + 0\mathbf{e}_2 + c_1\mathbf{e}_3$, where $\mathbf{w} = (a_1, c_1, 0, 0)^T \in \mathcal{X}$. For another solution $\mathbf{x}_2 = (a_1, c_1, c_1, b_2)^T \in \mathcal{X}$, since $\mathbf{w}_0 = (a_1, c_1, 0, b_2)^T$, and $\mathbf{x}_2 \geq \mathbf{w}_0$, we have $\mathbf{x}_2 = \mathbf{w}_0 \vee \mathbf{y}$, where $\mathbf{y} = (0, 0, c_1, 0)^T \in \mathcal{Y}$, that is, $\mathbf{x}_2 = \mathbf{w}_0 + 0\mathbf{e}_1 + 0\mathbf{e}_2 + c_1\mathbf{e}_3$. And these also demonstrate the description in Remark 4.1.

Now we consider a special case of Eq.(1), that is, $|I| = 1$, and assume that the equation is

$$\bigvee_{j \in J} (a_j \wedge x_j) = b, \quad (3)$$

and its associated equation is

$$\bigvee_{j \in J} (a_j \wedge x_j) = 0. \quad (4)$$

Obviously, the greatest solution of Eq.(4) is $\mathbf{y}^* = (a'_j)_{j \in J}^T$, and $\mathbf{e}_k = (y_j)_{j \in J}^T$, where $y_k = a_k$ and $y_j = 0$ if $j \neq k$.

There are some relationships between the minimal solution of Eq.(3) and the minimal decomposition of b .

Theorem 4.8 Let $J' \subseteq J$. If $M = \{s_j \in L : s_j \leq a_j \text{ for each } j \in J'\} \in \mathcal{M}_b$, then $\mathbf{x}_* = (x_{j*})_{j \in J}^T \in \mathcal{X}^0$, where $x_{j*} = s_j$ if $j \in J'$ and $x_{j*} = 0$ if $j \in J - J'$.

Proof. It is clear that $\mathbf{x}_* = (x_{j*})_{j \in J}^T \in \mathcal{X}$. If $\mathbf{x} = (x_j)_{j \in J}^T \in \mathcal{X}$ and $\mathbf{x} \leq \mathbf{x}_*$, then $x_j \leq x_{j*}$ for all $j \in J$. This implies that $x_j \leq x_{j*}$ if $j \in J'$ and $x_j = x_{j*} = 0$ if $j \in J - J'$. Thus $b = \bigvee_{j \in J'} x_j$. This last formula combined with the fact $M = \{s_j \in L : j \in J'\} \in \mathcal{M}_b$ yields that $x_j = s_j = x_{j*}$ for all $j \in J$. Consequently, $\mathbf{x}_* = (x_{j*})_{j \in J}^T \in \mathcal{X}^0$. ■

Theorem 4.9 If $\mathcal{X} \neq \emptyset$, then for any $\mathbf{x} \in \mathcal{X}$, there is some $\mathbf{x}_* \in \mathcal{X}^0$ such that $\mathbf{x}_* \leq \mathbf{x}$.

Proof. Suppose that $\mathbf{x} = (x_j)_{j \in J}^T \in \mathcal{X}$. Then $\bigvee_{j \in J} (a_j \wedge x_j) = b$. In this representation, Removing

the redundant elements, we can get a irredundant decomposition $b = \bigvee_{j \in J'} (a_j \wedge x_j)$, where $J' \subseteq J$ and $b \neq \bigvee_{j \in J', j \neq k} (a_j \wedge x_j)$ for each $k \in J'$. By Theorem 3.3, there exists $M = \{s_j \in L : j \in J'\} \in \mathcal{M}_b$ such that $s_j \leq a_j \wedge x_j$ for all $j \in J'$. Define $\mathbf{x}_* = (x_{j*})_{j \in J}^T$ with $x_{j*} = a_j \wedge x_j$ if $j \in J'$ and $x_{j*} = 0$ if $j \in J - J'$. It is clear that $\mathbf{x}_* \leq \mathbf{x}$ and $\mathbf{x}_* \in \mathcal{X}^0$. ■

According to Theorem 4.9, if $\mathcal{X} \neq \emptyset$, then $\mathcal{X}^0 \neq \emptyset$. Now let $\mathbf{x}_0 = \bigvee_{\mathbf{x} \in \mathcal{X}^0} \mathbf{x}$. Then we have the following statement:

Theorem 4.10 $\mathbf{x}_0 = (a_j \wedge b)_{j \in J}$ and $\mathbf{y}^* \vee \mathbf{x}_0 = \mathbf{x}^*$.

Proof. Let $\mathbf{x}_0 = (x_{j0})_{j \in J}^T$. By Lemma 2.6, for any $\mathbf{x}_* = (x_{j*})_{j \in J}^T \in \mathcal{X}^0$, we have $x_{j*} \leq a_j \wedge b$ for all $j \in J$. Then $x_{j0} \leq a_j \wedge b$ for all $j \in J$. In the following, we show that for any $k \in J$, there is some $\mathbf{x}_* = (x_{j*})_{j \in J}^T \in \mathcal{X}^0$ such that $x_{k*} = a_k \wedge b$. Obviously, if $a_k \wedge b = 0$, then $x_{k*} = 0$ for any $\mathbf{x}_* = (x_{j*})_{j \in J}^T \in \mathcal{X}^0$. If $a_k \wedge b = b$, that is, $a_k \geq b$, then $\mathbf{x}_* = (x_{j*})_{j \in J}^T \in \mathcal{X}^0$, where $x_{k*} = b$ and $x_{j*} = 0$ with $j \neq k$. If $0 < a_k \wedge b < b$, then $a_k \wedge b \in D(b)$. Since $b = \bigvee_{j \in J} (a_j \wedge b) = (a_k \wedge b) \vee \bigvee_{j \in J, j \neq k} (a_j \wedge b)$, by Theorem 3.4 it follows that there is a minimal decomposition $b = (a_k \wedge b) \vee \bigvee_{j \in J_k} u_j$, where $J_k \subseteq J - \{k\}$ and $u_j \leq a_j \wedge b$ for all $j \in J_k$. Define $\mathbf{x}_* = (x_{j*})_{j \in J}^T$ with $x_{k*} = a_k \wedge b$, $x_{j*} = u_j$ if $j \in J_k$ and $x_{j*} = 0$ if $j \in J - (J_k \cup \{k\})$. Then $\mathbf{x}_* \in \mathcal{X}^0$ by Theorem 4.8. With the above discussions, we have $\mathbf{x}_0 = (a_j \wedge b)_{j \in J}^T$. Now for any $j \in J$, $x_j^* = a'_j \vee b$ and $y_j^* = a'_j$, it follows that for all $j \in J$, $a'_j \vee (a_j \wedge b) = (a'_j \vee a_j) \wedge (a'_j \vee b) = 1 \wedge (a'_j \vee b) = a'_j \vee b = x_j^*$. Consequently, $\mathbf{y}^* \vee \mathbf{x}_0 = \mathbf{x}^*$. ■

Theorem 4.11 Let $\mathbf{x} \in \mathcal{X}$, and let $\mathbf{x}_1 = \bigvee_{\mathbf{x}_* \in \mathcal{X}^0, \mathbf{x}_* \leq \mathbf{x}} \mathbf{x}_*$. Then there is some element $\mathbf{y} \in \mathcal{Y}$ such that $\mathbf{x} = \mathbf{x}_1 \vee \mathbf{y}$.

Proof. Let $\mathbf{x} = (x_j)_{j \in J}^T \in \mathcal{X}$ and $\mathbf{x}_1 = (x_{j1})_{j \in J}^T$. Then $x_{j1} \leq a_j$ and $x_{j1} \leq x_j$, that is, $x_{j1} \leq a_j \wedge x_j$. In the following, we show that $x_{j1} = a_j \wedge x_j$ for all $j \in J$. It is clear that we only show that for any $k \in J$, there is some $\mathbf{x}_* = (x_{j*})_{j \in J}^T \in \mathcal{X}^0$ such that $\mathbf{x}_* \leq \mathbf{x}$ and $x_{k*} = a_k \wedge x_k$. Obviously, if $a_k \wedge x_k = 0$, then $x_{k*} = 0$ for any $\mathbf{x}_* = (x_{j*})_{j \in J}^T \in \mathcal{X}^0$ with $\mathbf{x}_* \leq \mathbf{x}$. If $a_k \wedge x_k = b$, that is, $a_k \geq b$, then $\mathbf{x}_* = (x_{j*})_{j \in J}^T \in \mathcal{X}^0$, where $x_{k*} = b$ and $x_{j*} = 0$ with $j \neq k$. If $0 < a_k \wedge x_k < b$, then $a_k \wedge x_k \in D(b)$. Similar to the proof of Theorem 4.10, there is a minimal solution $\mathbf{x}_* = (x_{j*})_{j \in J}^T \in \mathcal{X}^0$ such that $\mathbf{x}_* \leq \mathbf{x}$ and $x_{k*} = a_k \wedge x_k$. With the above discussions, we have $\mathbf{x}_1 = (a_j \wedge x_j)_{j \in J}^T$. Now let $\mathbf{y} = \mathbf{y}^* \wedge \mathbf{x}$, then $\mathbf{y} \in \mathcal{Y}$. Furthermore, for every $j \in J$, $(a_j \wedge x_j) \vee (a'_j \wedge x_j) = (a_j \vee a'_j) \wedge x_j = 1 \wedge x_j = x_j$, that is, $\mathbf{x} = \mathbf{x}_1 \vee \mathbf{y}$. ■

In particular, we have the following statements.

Corollary 4.1 For any $\mathbf{x} \in \mathcal{X}$, if $\mathbf{x}_0 \leq \mathbf{x} \leq \mathbf{x}^*$, then there is a solution $\mathbf{y} \in \mathcal{Y}$ such that $\mathbf{x} = \mathbf{x}_0 \vee \mathbf{y}$.

Corollary 4.2 *If \mathcal{X} has the least solution \mathbf{x}_* , then for any $\mathbf{x} \in \mathcal{X}$, there is a solution $\mathbf{y} \in \mathcal{Y}$ such that $\mathbf{x} = \mathbf{x}_* \vee \mathbf{y}$.*

Remark 4.3 *According to Theorem 4.11, we have that for any $\mathbf{x} = (x_j)_{j \in J}^T \in \mathcal{X}$ and $\mathbf{x}_1 \leq \mathbf{x} \leq \mathbf{x}^*$, \mathbf{x} can be represented by a linear combination of a special solution \mathbf{x}_1 of \mathcal{X} and some certain solutions of \mathcal{Y} , that is, $\mathbf{x} = \mathbf{x}_1 + y_1 \mathbf{e}_1 + y_2 \mathbf{e}_2 + \dots + y_n \mathbf{e}_n$, where $\mathbf{y} = \mathbf{y}^* \wedge \mathbf{x} = (y_j)_{j \in J}^T \in \mathcal{Y}$.*

Example 4.2 *Let L be the lattice from Example 4.1. Consider equation $(c_1 \wedge x_1) \vee (c_2 \wedge x_2) \vee (b_4 \wedge x_3) \vee (a_1 \wedge x_4) \vee (a_2 \wedge x_5) = b_5$. Obviously, $\mathcal{X} \neq \emptyset$, and $\mathbf{x}^* = (a_4, 1, a_2, b_5, a_4)^T$, $\mathcal{X}^0 = \{\mathbf{x}_{*1} = (0, 0, 0, 0, b_5)^T, \mathbf{x}_{*2} = (0, c_2, 0, 0, c_4)^T, \mathbf{x}_{*3} = (0, 0, c_2, 0, c_4)^T, \mathbf{x}_{*4} = (0, 0, 0, c_2, c_4)^T\}$. The associated equation is $(c_1 \wedge x_1) \vee (c_2 \wedge x_2) \vee (b_4 \wedge x_3) \vee (a_1 \wedge x_4) \vee (a_2 \wedge x_5) = 0$, and $\mathbf{y}^* = (a_4, a_3, b_3, c_4, c_3)^T$, $\mathbf{e}_1 = (a_4, 0, 0, 0, 0)^T$, $\mathbf{e}_2 = (0, a_3, 0, 0, 0)^T$, $\mathbf{e}_3 = (0, 0, b_3, 0, 0)^T$, $\mathbf{e}_4 = (0, 0, 0, c_4, 0)^T$, $\mathbf{e}_5 = (0, 0, 0, 0, c_3)^T$. Consider a solution $\mathbf{x} = (b_4, a_1, b_1, 0, b_6)^T \in \mathcal{X}$, since $\mathbf{x} \geq \mathbf{x}_{*2} \vee \mathbf{x}_{*3} = (0, c_2, c_2, 0, c_4) = \mathbf{x}_1$, $\mathbf{y}^* \wedge \mathbf{x} = (b_4, b_2, c_1, 0, c_3)^T$, it follows that $\mathbf{x} = \mathbf{x}_1 + b_4 \mathbf{e}_1 + b_2 \mathbf{e}_2 + c_1 \mathbf{e}_3 + 0 \mathbf{e}_4 + c_3 \mathbf{e}_5$.*

5. Conclusions

In this contribution, we investigated the structure of the solution set of Eq.(1) over complete Boolean algebras similar to linear algebraic systems, and showed that each solution can be represented by a linear combination of a special solution of Eq.(1) and some certain solutions of the homogeneous equations associated with Eq. (1). In Eq.(3), the special solution can be determined by the minimal solutions. Unfortunately, without additional conditions, for Eq.(1), the special solution is not unique.

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